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Secure evaluation of polynomial using privacy ring homomorphisms

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Abstract: Method of secure evaluation of polynomial $y=F(x_1, \dots, x_k)$ over some rings on untrusted computer is proposed. Two models of untrusted computer are considered: passive and active. In passive model untrusted computer correctly computes polynomial F and tries to know secret input (x_1, \dots, x_k) and output y . In active model untrusted computer tries to know input and output and tries to change correct output y so that this change cannot be determined. Secure computation is proposed by using one-time privacy ring homomorphism $\mathbb{Z}/n\mathbb{Z} \rightarrow \mathbb{Z}/n\mathbb{Z}[z]/(f(z))$, $n = pq$, generated by trusted computer. In the case of active model secret check point $v = F(u_1, \dots, u_k)$ is used. Trusted computer generates polynomial $f(z)=(z-t)(z+t)$, t in $\mathbb{Z}/n\mathbb{Z}$, and input $X_i(z)$ in $\mathbb{Z}/n\mathbb{Z}[z]/(f(z))$ such that $X_i(t)=x_i \pmod{n}$ for passive model, and $f(z)=(z-t_1)(z-t_2)(z-t_3)$, t_i in $\mathbb{Z}/n\mathbb{Z}$ and input $X_i(z)$ in $\mathbb{Z}/n\mathbb{Z}[z]/(f(z))$ such that $X_i(t_1)=x_i \pmod{n}$, $X_i(t_2)=u_i \pmod{n}$ for active model. Untrusted computer computes function $Y(z) = F(X_1(z), \dots, X_k(z))$ in the ring $\mathbb{Z}/n\mathbb{Z}[z]/(f(z))$. For passive model trusted computer determines secret output $y=Y(t) \pmod{n}$. For active model trusted computer checks that $Y(t_2)=v \pmod{n}$, then determines correct output $y=Y(t_1) \pmod{n}$.

Category / Keywords: cryptographic protocols / elliptic curve cryptosystem, factoring, public-key cryptography

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