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Non-malleable Codes from Additive Combinatorics

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Abstract: Non-malleable codes provide a useful and meaningful security guarantee in situations where traditional error-correction (and even error-detection) is impossible; for example, when the attacker can completely overwrite the encoded message. Informally, a code is non-malleable if the message contained in a modified codeword is either the original message, or a completely unrelated value. Although such codes do not exist if the family of "tampering functions" \mathcal{C} is completely unrestricted, they are known to exist for many broad tampering families \mathcal{C} . One such natural family is the family of tampering functions in the so called $\{\text{em split-state}\}$ model. Here the message m is encoded into two shares L and R , and the attacker is allowed to arbitrarily tamper with L and R $\{\text{em individually}\}$. The split-state tampering arises in many realistic applications, such as the design of non-malleable secret sharing schemes, motivating the question of designing efficient non-malleable codes in this model.

Prior to this work, non-malleable codes in the split-state model received considerable attention in the literature, but were either (1) constructed in the random oracle model [DPW10], or (2) relied on advanced cryptographic assumptions (such as non-interactive zero-knowledge proofs and leakage-resilient encryption) [LL12], or (3) could only encode 1-bit messages [DKO13]. As our main result, we build the first efficient, multi-bit, information-theoretically-secure non-malleable code in the split-state model.

The heart of our construction uses the following new property of the inner-product function $\langle L, R \rangle$ over the vector space F^n (for any finite field F and large enough dimension n): if L and R are uniformly random over F^n , and $f, g: F^n \rightarrow F^n$ are two arbitrary functions on L and R , the joint distribution $(\langle L, R \rangle, \langle f(L), g(R) \rangle)$ is "close" to the convex combination of "affine distributions" $\{(U, cU+d) \mid c, d \in F\}$, where U is uniformly random in F . In turn, the proof of this surprising property of the inner product function critically relies on some results from additive combinatorics, including the so called $\{\text{em Quasi-polynomial Freiman-Ruzsa Theorem}\}$ (which was recently established by Sanders [San12] as a step towards resolving the Polynomial Freiman-Ruzsa conjecture [Gre05]).

Category / Keywords: applications / Non malleable codes, Combinatorics

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