Performance Analysis of Some Password Hashing Schemes

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Abstract

In this work we have analyzed some password hashing schemes for performance under various settings of time and memory complexities. We have attempted to benchmark the said algorithms at similar levels of memory consumption. Given the wide variations in security margins of the algorithms and incompatibility of memory and time cost settings, we have attempted to be as fair as possible in choosing the various parameters while executing the benchmarks.

Keywords: Password hashing, benchmark, PHC

1 Benchmarking setup

In order to get consistent results for the different algorithms we performed all the test on a single machine with the code compiled by the same compiler. The details are as follows:

- CPU: Intel Core i7 4770 (Turbo Boost: ON) Working at 3.9 GHz
- RAM: Double Channel DDR3 16 GB (2400 MHz)
- Compiler: gcc / g++ v4.9.2 (-march=native and -O3 flags were set if not already in the makefiles). This would cause the compiler to use the AVX-2 instructions.
- **OS:** UBUNTU 14.04.1, on HYPER-V, on Windows-8.1 with 8 GB allocated RAM to the VM. We also performed benchmarks on native Linux OS to make sure that the virtualization does not cause any changes in the results.

2 Performance Analysis

In this section we provide the benchmarking results and details of the setup and considerations. For consistency we used single threaded versions of the algorithms.

All the experiments were run at-least 5 times and the average timings were taken.

2.1 Catena v3 - Butterfly and Dragonfly [6]

The Catena v3 document provides the two new variants called Dragonfly and Butterfly based on the instantiation of Catena-BRG and Catena-DBG. Earlier versions were significantly slower than the current version with the single round Blake2b-1 function.

1	Cater	ıa	v3 Butterfly	7	 (at di	fferer	nt	values of la	aml	oda with BLAKI	E2b-1)
1	t_cost	•		-					-	•	2048 MB (m_cost=24)
 	lambda=1	•		•			•		•	13.4 sec 74.0 MB/s	
 	lambda=2	•		•			•		•	26.9 sec 38.0 MB/s	
 	lambda=3	•		-					-	40.0 sec 25.3 MB/s	

Table 1: Catena v3 - Butterfly

The Catena-Dragonfly is much faster than Catena-Butterfly, but, the Dragonfly version is shown not to be memory-hard in [4], [2].

The latest version of code at the time of benchmark was cloned from 'https://github.com/cforler/catena/' (75 commits).

Optimized SSE implementation was used for both the benchmarks.

Catena v	3 Dragonfly	(at dif	ferent	values of	lambda	with I	BLAKE2b-1)	
t_cost	128 MB	256 MI	3 I	512 MB	l 1024	MB	2048 MB	1
lambda=1	0.237 sec 539 MB/s						•	
lambda=2	0.29 sec 437 MB/s	•	•		•		•	-
lambda=3	0.46 sec 273 MB/s	•	•		•		•	-
lambda=4	0.534 sec 239 MB/s	•	•		•		•	-

Table 2: Catena Dragonfly

One of the reasons for the slow nature of the Butterfly version is the need for $2 \cdot g$ rows for processing. This property of Catena-DBG, combined with the relatively small read-writes to the RAM makes the overall structure significantly slow. Even the fastest version of Catena-Butterfly-Blake2b-1 can only achieve overall memory hashing speed of around 80 MiB/s.

The Catena-Dragonfly (Table 2) is much faster due to the significantly reduced number of rounds as compared to Catena-Butterfly (Table 1). Also, the way it operates, every node in the Catena-BRG graph has dependency on two previous ancestors as opposed to three of Catena-DBG. This leads to reduced number of random memory accesses and faster speeds.

2.2 Gambit[10]

For the benchmarking of Gambit we used v1 of the source code from [1]. No optimized version of the code was available, and we used the reference implementation for this analysis. The speeds are particularly slow due to the slow performance of Keccak sponge in software and small memory access chunks. One peculiarity of Gambit is that the Time Cost and Memory Cost is bound by the assertion $(cost_m \times 2 \le cost_t \times (r/8))$ and r = 136 for Gambit - 256. For the benchmark we set t_cost to the lowest possible value for a required m_cost for fixed memory consumption and defined it as t = 1, for higher values of t we doubled the t_cost for every subsequent values of t. This was done to have a consistent range of possible speeds with increasing t. Results are shown in Table 3.

l Ga	 nbit-v1	Memory proces	ssing rate of	Gambit-256
t_cos	t 128 MB	256 MB	512 MB	1024 MB
t=1 15978	•	ec 2.47 se MB/s 103 MB,		ec 9.57 sec MB/s 106.9 MB/s
t=2 31956	-	•	·-	sec 19.62 sec MB/s 52.18 MB/s
t=3 47934				sec 28.75 sec IB/s 35.6 MB/s
t=4 63913	•	ec 9.82 se MB/s 26 MB/s		sec 38.14 sec IB/s 26.84 MB/s

Note: t_cost values are for 128 MB.

Table 3: Gambit v1

2.3 Lyra2 - v3[7]

For benchmarking Lyra-2 we used the v3 code from [1]. We used the SSE version of Lyra2-v3 with nPARALLEL=1 for consistency. We noticed that the speed of Lyra-2 with multiple threads is faster than the single thread ones (as expected), but for consistency, we choose to use the single threaded version. The default Makefile was used with linux-x86-64-sse, with nThreads=1. This would result in fast AVX implementation being used with Blake2b as sponge function. We did receive a warning for large-memory-usage for the 2 GiB test, but, the timings are as expected.

Results are shown in Table 4.

-	L	/ra	a2 - 1	73 (a	at	diffe	ent	va.	lues o	of ta	ano	d p=1))							1
	t_cost		200	MB	1	400 ME	3		800	MB		1024	MB		1600 M	iB	1	2048	MB	1
 	t=1	•						-			-				0.98 se		•			-
 	t=2	•						-			-				1.47 se 1088 ME		•			-
 	t=3	•						-			-				1.91 se 837 MB/		•			-
 	t=4	 													2.38 se 669 ME					

Table 4: Lyra2 - v3

2.4 Rig v2[3]

For this work we used the latest version of Rig from 'https://github.com/arpanj/Rig'. We used the optimized implementation with the Blake2b round using AVX-2.

 I	Ι	RIC	i v	2	 (Bl	akeEx	par	 nd, Bl	akePe	rm	 , Blak	 e2b)		 AVX-2	x86-6	 64			
m	ı =	=>		_	13 (1	28 M)	 	14 (2	56 M)		15 (5:	12 M)		16 (10	024 M)		17 (20	048 M)	
n	. =	= 1					-							0.519 1973		-			
n	. =	= 2					-							0.718 1425		-			
n	. =	= 3	•				•			•			•	0.947 1081		•			
n	. =	= 4					-							1.168 876		-			

Table 5: RIG v2

All default settings were used as described in the code and Makefile. One source code improvement was the removal of writing of the data back to the memory in the last row, this change resulted in around five percent improvement in overall performance for small values of N. Results are shown in Table 5.

2.5 Scrypt[8]

Scrypt is the first memory-hard algorithm for password-hashing. There are several implementations of Scrypt available, we used one of the fastest variants of the implementation by @floodyberry in this work. Table 6 shows the results of the AVX2 implementation with Blake2b and Salsa64/8.

Scrypt @floodyberry's https://github.com/floodyberry/scrypt-jane

I	Scrypt (AVX2	2, B	lake2b	o, Salsa	64/8,	x86-64)	I
 	Memory (MB)		Time	(second)		Speed (MB/s)	
ı	128	I		0.076	1	1684	ı
1	256	- 1		0.162	1	1580	-
1	512	- 1		0.332	- 1	1542	-
1	1024	-		0.7	- 1	1530	-1

Table 6: scrypt: floodyberry/scrypt-jane

2.6 TwoCats[5]

TwoCats is one of the fastest and one of the most complex entries of [1]. It is highly optimized to use the CPU and memory subsystem to the full extent by having several modes and multi-threading support. For the purpose of this analysis, we used the single threaded mode of TwoCats with Blake2b compiled with AVX2 support. Defaults were used among them MULTIPLIES=2, LANES=8, BLOCKSIZE=16384 and PARALLELISM=1.

The t_cost parameter of TwoCats is quite sensitive as it increases the iteration count of the number of small writes in cache using 2^{t_cost} .

The results are shown in Table 7.

2.7 yescrypt [9]

For benchmarking yescrypt, we used version 0.7.1 of the code from [1]. Yescrypt is another fast and complex submission to the PHC. There are several modes and settings available. For this work we used the default configuration (r=8, p=1 and YESCRYPT_RW=1). 64 bit version was used with -march=native in gcc (which essentially would have enabled AVX2 intrinsic support). The code however uses 128 bit intrinsics, so SSE4.1 should be enough to compile and enable SIMD optimizations.

The results are shown in Table 8.

1]	 Γw	Cats	v0 (at	5 6	differ	ent v	alı	es of	t wit	ch	Blake:	2b)				 I
I	t_cost	1	128	MB	1	256	MB	1	512	MB	1	1024	MB	I	2048	MB	1
1	t=0				- :										1.025 1997		
	t=1										•			-	1.842 1111		
 	t=2										•			-	3.514 582		
	t=3														6.810 300		
	t=4										•			-	13.49 152		

Table 7: TwoCats v0

1			yes	crypt	v1	L (at \$	SHA256	3,	r=8,]	p=1 ai	nd	YESCR	YPT_R	√=:	1)		
1	t_cost	1	128	MB	1	256	MB	1	512	MB		1024	MB		2048	MB	
1	t=0	1			-										1.46 1402		
1	t=1	-			-						-			-	1.782 1149		
 	t=2														2.1 975		
 	t=3	-			-						-			-	3.15 650		
 	t=4														4.146 493		

Table 8: Yescrypt v1

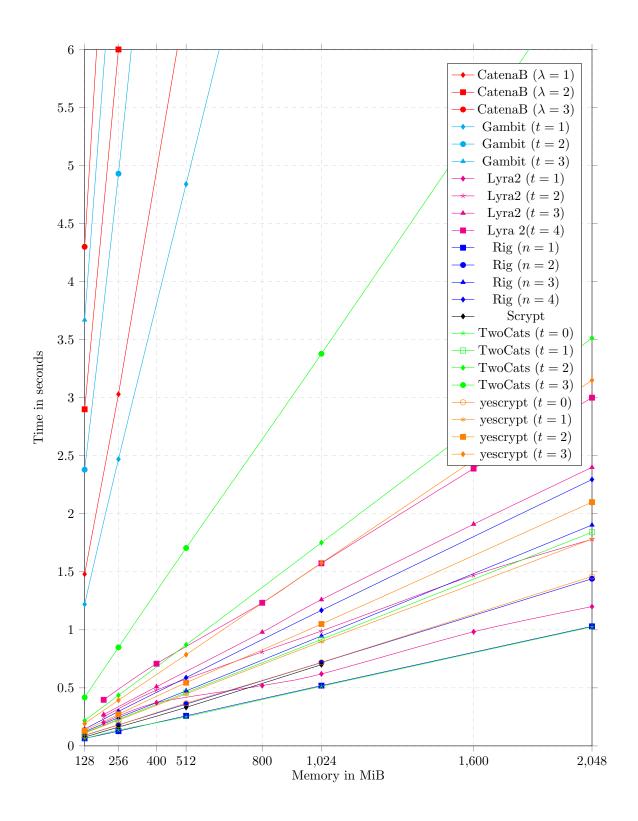


Figure 1: Performance: Memory vs. Time of some memory-hard PHC candidates

3 Conclusions

The performance graph in figure 1 shows the execution time vs. memory for all the memory-hard algorithms benchmarked. It is clear that Gambit (reference code) and Catena-Butterfly are among the slowest and take significant amount of time in hashing passwords with moderate to large amounts of memory. The performance of Gambit may be improved using a better implementation, but, the performance of Catena is unlikely to significantly improve even with native assembly implementation.

As noted before, the time cost of TwoCats is very sensitive, it may be changed with some minor tweaks to allow for better tradeoff control.

Lyra2, Rigv2, TwoCats and yescrypt provide good performance in a wide range of use cases. No, attacks are currently known against them which reduce the claimed TMTO defense.

As far as side-channels are concerned, Catena, Gambit and Rig are fully resistant; Lyra2 and TwoCats are partially resistant whereas Scrypt and yescrypt are not resistant.

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