

Classification of Elliptic/hyperelliptic Curves with Weak Coverings against GHS Attack without Isogeny Condition

Tsutomu Iijima ^{*} Fumiyuki Momose [†] Jinhui Chao [‡]

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Abstract

The GHS attack is known as a method to map the discrete logarithm problem (DLP) in the Jacobian of a curve C_0 defined over the d degree extension k_d of a finite field k to the DLP in the Jacobian of a new curve C over k which is a covering curve of C_0 .

Recently, classification and density analysis were shown for all elliptic and hyperelliptic curves C_0/k_d of genus 2, 3 which possess $(2, \dots, 2)$ covering C/k of \mathbb{P}^1 under the isogeny condition (i.e. when $g(C) = d \cdot g(C_0)$). In this paper, we show a complete classification of small genus hyperelliptic curves C_0/k_d which possesses $(2, \dots, 2)$ covering C over k without the isogeny condition. Our main approach is to use representation of the extension of $Gal(k_d/k)$ acting on $cov(C/\mathbb{P}^1)$. In the classification we restricted the group order or key-length of the DLP to certain range reasonable in cryptographic application. Explicit defining equations of such curves and the existence of a model of C over k are also presented.

Keywords : Weil descent attack, GHS attack, Elliptic curve cryptosystems, Hyperelliptic curve cryptosystems, Index calculus, Galois representation

1 Introduction

Let $k_d := \mathbb{F}_{q^d}$, $k := \mathbb{F}_q$ ($d > 1$), q be a power of a prime number.

Weil descent was firstly introduced by Frey [8] to elliptic curve cryptosystems. This idea is developed into the well-known GHS attack in [12]. This attack maps the discrete logarithm problem (DLP) in the Jacobian of a curve

^{*}Koden Electronics Co.,Ltd., 2-13-24 Tamagawa, Ota-ku, Tokyo, 146-0095 Japan

[†]Department of Mathematics, Chuo University, 1-13-27 Kasuga, Bunkyo-ku, Tokyo, 112-8551 Japan

[‡]Department of Information and System Engineering, Chuo University, 1-13-27 Kasuga, Bunkyo-ku, Tokyo, 112-8551 Japan

C_0 defined over the d degree extension field k_d of the finite field k to the DLP in the Jacobian of a curve C over k by a conorm-norm map. The GHS attack is further extended and analyzed in [2][4][9] [10][15][16][17][20][25][26], and is conceptually generalized to the cover attack [6]. The cover attack maps the DLP in the Jacobian of a curve C_0/k_d to the DLP in the Jacobian of a covering curve C/k of C_0 when a covering map or a non-constant morphism between C_0 and C exists.

If the DLP in the Jacobian of C_0 can be solved more efficiently in the Jacobian of C , we call C_0 a weak curve or say that it has weak covering C against GHS or cover attack. Thus, it is important and interesting to know what kind of curves C_0 have such coverings C , how many are they, etc..

It is known that the most efficient attack to DLP in the Jacobian of algebraic curve based systems is the index calculus algorithms. In [11], Gaudry first proposed his variant of the Adleman-DeMarrais-Huang algorithm [1] to attack hyperelliptic curve discrete logarithm problems, which is faster than Pollard's rho algorithm when the genus is larger than 4 but becomes impractical for large genera. Recently, a single-large-prime variation [27] and a double-large-prime variation [13][24] are proposed. These variations can be applied in the GHS attack if the curve C/k is a hyperelliptic curve of $g(C) \geq 3$. The complexity of these double-large-prime algorithms are $\tilde{O}(q^{2-2/g})$. On the other hand, when C/k is a non-hyperelliptic curve, Diem's recent proposal of a double-large-prime variation [5] can be applied with complexity of $\tilde{O}(q^{2-2/(g-1)})$. This algorithm is not only faster than Pollard's rho algorithm but also the fastest attack algorithm to curve based cryptosystems at present.

Recently, a thorough security analysis of elliptic and hyperelliptic curves C_0/k_d with weak covering C/k is shown in [3][21][22][23] under the following isogeny condition. Assuming that there exists a covering curve C/k of C_0/k_d ,

$$\exists \pi/k_d : C \longrightarrow C_0 \tag{1}$$

such that for

$$\pi_* : J(C) \longrightarrow J(C_0), \tag{2}$$

$$Res(\pi_*) : J(C) \longrightarrow Res_{k_d/k} J(C_0) \tag{3}$$

is an isogeny, here $J(C)$ is the Jacobian variety of C and $Res_{k_d/k} J(C_0)$ is its Weil restriction. Then $g(C) = d \cdot g(C_0)$.

Under this isogeny condition, C_0/k_d which possesses covering curves C/k as $(2, \dots, 2)$ covering of \mathbb{P}^1 are classified for hyperelliptic curves of genus 1,2,3 in [3][14][21][22][23]. Density and defining equations are also presented for these curves. Further in [18], when $g(C) = d \cdot g(C_0) + e$, ($e > 0, d = 2, 3, 4$) for $g(C_0) = 1, 2, 3$ hyperelliptic curves in the cryptographic applications, certain classes of curves C_0/k_d which have weak coverings C/k were showed.

In this paper, we show a complete classification of hyperelliptic curves C_0/k_d of genus 1,2,3 with $(2, \dots, 2)$ covering C/k without isogeny condition. In particular, we assume that $g(C) = d \cdot g(C_0) + e, e > 0$. The classification is then restricted to a certain range of the group order or key-length reasonable for cryptographic applications. Our approach for the classification is a representation theoretical one, to investigate action of the extension of $Gal(k_d/k)$ on $cov(C/\mathbb{P}^1)$. We also present defining equations of these curves and existential conditions of a model of C over k explicitly.

2 GHS attack and $(2, \dots, 2)$ covering

Firstly, we summarize the GHS attack and the cover attack. Let $k_d(C_0)$ be the function field of a curve C_0/k_d , $Cl^0(k_d(C_0))$ the class group of the degree 0 divisors of $k_d(C_0)$, $\sigma_{k_d/k}$ the Frobenius automorphism of k_d over k . Assume $\sigma_{k_d/k}$ is extended to an automorphism σ of order d in the separable closure of $k_d(x)$. The Galois closure of $k_d(C_0)/k(x)$ is $F' := k_d(C_0) \cdot \sigma(k_d(C_0)) \cdots \sigma^{d-1}(k_d(C_0))$ and the fixed field of F' by the automorphism σ is $F := \{\alpha \in F' \mid \sigma(\alpha) = \alpha\}$. The DLP in $Cl^0(k_d(C_0))$ is mapped to the DLP in $Cl^0(F)$ using the following composition of conorm and norm maps:

$$N_{F'/F} \circ Con_{F'/k_d(C_0)} : Cl^0(k_d(C_0)) \longrightarrow Cl^0(F).$$

This map is called the conorm-norm homomorphism in the original GHS paper on the elliptic curve case [12].

This attack has been extended to wider classes of curves [2][4][9][10][15][16][17][25][26]. The GHS attack is conceptually generalized to the cover attack by Frey and Diem [6]. When there exist an algebraic curve C/k and a covering $\pi/k_d : C \longrightarrow C_0$, the DLP in $J(C_0)(k_d)$ can be mapped to the DLP in $J(C)(k)$ by a pullback-norm map.

$$\begin{array}{ccc} J(C)(k_d) & \xleftarrow{\pi^*} & J(C_0)(k_d) \\ N \downarrow & \swarrow N \circ \pi^* & \\ J(C)(k) & & \end{array}$$

Hereafter, let q be a power of an odd prime. Assume C_0 is a $g(C_0) \in \{1, 2, 3\}$ hyperelliptic curve given by

$$C_0/k_d : y^2 = f(x). \tag{4}$$

Then we have a tower of extensions of function fields such that $k_d(x, y, \sigma^1 y, \dots, \sigma^{n-1} y)$ / $k_d(x)$ ($n \leq d$) is a $\overbrace{(2, \dots, 2)}^n$ type extension. Here, a $\overbrace{(2, \dots, 2)}^n$ covering is

defined as a covering $\pi/k_d : C \longrightarrow \mathbb{P}^1$

$$C \longrightarrow \underbrace{C_0 \longrightarrow \mathbb{P}^1(x)}_2 \quad \overbrace{(2, \dots, 2)}^n \quad (5)$$

such that $\text{cov}(C/\mathbb{P}^1) \simeq \mathbb{F}_2^n$, here $\text{cov}(C/\mathbb{P}^1) := \text{Gal}(k_d(C)/k_d(x))$.

3 Representation of $\text{Gal}(k_d/k)$ on $\text{cov}(C/\mathbb{P}^1)$

Next, we consider the Galois group $\text{Gal}(k_d/k)$ acting on the covering group $\text{cov}(C/\mathbb{P}^1) \simeq \mathbb{F}_2^n$.

$$\text{Gal}(k_d/k) \curvearrowright \text{cov}(C/\mathbb{P}^1) \simeq \mathbb{F}_2^n \quad (6)$$

Then one has a map onto $\text{Aut}(\text{cov}(C/\mathbb{P}^1))$.

$$\xi : \text{Gal}(k_d/k) \hookrightarrow \text{Aut}(\text{cov}(C/\mathbb{P}^1)) \simeq \text{GL}_n(\mathbb{F}_2) \quad (7)$$

Then, the representation of σ for given n, d is as follows:

$$\sigma = \left(\begin{array}{cccc} \boxed{\spadesuit_1} & O & \cdots & O \\ O & \boxed{\spadesuit_2} & \ddots & \vdots \\ \vdots & \ddots & \ddots & O \\ O & \cdots & O & \boxed{\spadesuit_s} \end{array} \right) \left. \begin{array}{l} \} n_1 \\ \\ \\ \} n_s \end{array} \right\} n = \sum_{i=1}^s n_i \quad (8)$$

where the O is a zero matrix,

$$\boxed{\spadesuit_i} = \begin{pmatrix} \boxed{\star_i} & \boxed{\star_i} & \tilde{O} & \cdots \\ \tilde{O} & \boxed{\star_i} & \ddots & \ddots \\ \vdots & \ddots & \ddots & \boxed{\star_i} \\ \tilde{O} & \cdots & \tilde{O} & \boxed{\star_i} \end{pmatrix} \left. \begin{array}{l} 1 \\ \vdots \\ l_i \end{array} \right\} l_i \quad (9)$$

is an $n_i \times n_i$ matrix which has a form of an $l_i \times l_i$ block matrix. The sub-block $\boxed{\star_i}$ is an $n_i/l_i \times n_i/l_i$ matrix and \tilde{O} is an $n_i/l_i \times n_i/l_i$ zero matrix. Here, if $F_i(x) := (\text{the characteristic polynomial of } \boxed{\star_i})^{l_i}$, then $F(x) := \text{LCM}\{F_i(x)\}$ is the minimal polynomial of σ . Obviously, $F_i(\sigma) = 0$ and $F(\sigma) = 0$. When $d_i := \text{ord}(\boxed{\spadesuit_i})$, $d = \text{LCM}\{d_i\}$.

The examples of the representation of σ for given n and d are as follows:

Example 3.1. $n = 2, d = 2$

$$\sigma = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \quad (10)$$

$$F(\sigma) = (\sigma + 1)^2 = 0$$

Example 3.2. $n = 2, d = 3$

$$\sigma = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \quad (11)$$

$$F(\sigma) = \sigma^2 + \sigma + 1 = 0$$

Example 3.3. $n = 3, d = 3$

$$\sigma = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix} \quad (12)$$

$$F(\sigma) = (\sigma + 1)(\sigma^2 + \sigma + 1) = 0$$

Example 3.4. $n = 4, d = 6$

$$\sigma = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \text{ or } \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \quad (13)$$

$$F(\sigma) = (\sigma^2 + \sigma + 1)^2 = 0 \text{ or } F(\sigma) = (\sigma + 1)^2(\sigma^2 + \sigma + 1) = 0.$$

Notice that

$$\boxed{\star_1} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \text{ or } \boxed{\spadesuit_1} = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \boxed{\spadesuit_2} = \boxed{\star_2} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \text{ respectively.} \quad (14)$$

4 Upper bound of e in $g(C) = dg(C_0) + e$

From now, we consider the case of a hyperelliptic curve C_0/k_d for $g(C_0) \in \{1, 2, 3\}$ such that there is a covering $\pi/k_d : C \rightarrow C_0$ and the covering curve C/k has genus $g(C) = d \cdot g(C_0) + e$ ($e > 0$). Here, e can be regarded as the dimension of $\ker(\text{Res}(\pi_*))$. Firstly, for C_0 which are used in the cryptographic applications, we will estimate an upper bound of e for $g(C_0) \in \{1, 2, 3\}$. In algebraic curve based cryptosystems, the standard key length is above 160 bits at present. This means the size of the Jacobian of C_0/k_d is

$$q^{g(C_0)d} \geq 2^{160}. \quad (15)$$

We assume that the size of Jacobian of C/k is $q^{dg_0+e} \leq 2^a$.

Remark 4.1. *In this paper, we discuss within $a \leq 320$. However, the procedures in the section 4.3 can apply to any a such that $q^{dg_0+e} \leq 2^a$. Besides, we notice that Lemma 5.1 and the procedure in the section 5.2 are independent of choice of the range.*

4.1 Case $g(C_0) = 1$

Then, we have the following situation for $g_0 = 1$

$$\begin{cases} q^{d+e} \leq 2^a \\ 2^{160} \leq q^d. \end{cases} \quad (16)$$

Now, since $\frac{q^{d+e}}{q^d} = \frac{2^a}{2^{160}}$, $q^e \leq 2^{a-160}$. Consequently,

$$\log q^e \leq \log 2^{a-160}.$$

It follows that an upper bound of e is

$$e \leq \frac{(a-160)d}{160}. \quad (17)$$

When we assume $a \leq 320$, $e \leq d$ is obtained.

4.2 Case $g(C_0) = 2, 3$

Similarly, when $g(C_0) = 2$, assume that

$$\begin{cases} q^{2d+e} \leq 2^a \\ 2^{160} \leq q^{2d}. \end{cases} \quad (18)$$

Then $e \leq 2d$ if $a \leq 320$. When $g(C_0) = 3$, the double-large-prime algorithms have the cost of $\tilde{O}(q^{\frac{4}{3}d})$. Accordingly, the condition $q^{3d} \geq 2^{180}$ (i.e. $q^{\frac{4}{3}d} \geq 2^{80}$) should be adopted instead of $q^{3d} \geq 2^{160}$ ($q^{\frac{4}{3}d} \geq 2^{71.11\dots}$) to keep the same security level with $g_0 = 1, 2$ hyperelliptic curves (the costs of attack to each DLP are $q^{\frac{d}{2}} \geq 2^{80}$ for $g_0 = 1$, $q^d \geq 2^{80}$ for $g_0 = 2$ respectively). Thus, one can assume

$$\begin{cases} q^{3d+e} \leq 2^a \\ 2^{180} \leq q^{3d}. \end{cases} \quad (19)$$

Consequently, $e \leq \frac{7}{3}d$ if $a \leq 320$. In the next subsection, we enumerate the candidates of n, d, e, S within these bounds of e for $g(C_0) = 1, 2, 3$.

4.3 The candidates of (n, d, e, S)

Let S be the number of fixed points of C/\mathbb{P}^1 covering. By the Riemann-Hurwitz theorem, $2g-2 = 2^n(-2) + 2^{n-1}S$, then $S = 4 + \frac{dg_0+e-1}{2^{n-2}}$. Hereafter, we consider the following two types:

- Type (A) : $\exists d_i$ s.t. $d_i = d (= LCM\{d_i\})$
then, $S = 4 + \frac{dg_0+e-1}{2^{n-2}} \geq \max\{d, 2g_0 + 3\}$
- Type (B) : $d_i \neq d$ for $\forall d_i$
then, $S = 4 + \frac{dg_0+e-1}{2^{n-2}} \geq \max\{q(d), 2g_0 + 4\}$

here $q(d) := \sum p_i^{e_i}$ for $d = \prod p_i^{e_i}$ (p_i 's are distinct prime numbers). See the example 3.4 again. We notice the left and right matrices are a type (A) and a type (B) respectively.

4.3.1 Type (A)

- Case $g_0 = 1$:

From the above, $d + e - 1 \geq 2^{n-2}d - 2^n$ when $g_0 = 1$. Since we assume $0 < e \leq d$, $2d - 1 \geq d + e - 1 \geq 2^{n-2}d - 2^n$. Then $2^n - 1 \geq (2^{n-2} - 2)d$ ($n \geq 3$). Now, if $n > 3$,

$$(n \leq) d \leq 4 + \frac{7}{2^{n-2} - 2}. \quad (20)$$

Consequently, it follows that $n \geq 6$ is not within the candidates. From this result and the property of σ , the candidates of 4-triple (n, d, e, S) are: $(5, 5, 4, 5)$, $(4, 4, 1, 5)$, $(4, 5, 4, 6)$, $(4, 6, 3, 6)$, $(4, 7, 6, 7)$, $(3, 3, 2, 6)$, $(3, 4, 1, 6)$, $(3, 4, 3, 7)$, $(3, 7, 2, 8)$, $(3, 7, 4, 9)$, $(3, 7, 6, 10)$, $(2, 2, 1, 6)$, $(2, 2, 2, 7)$, $(2, 3, 1, 7)$, $(2, 3, 2, 8)$, $(2, 3, 3, 9)$.

- Case $g_0 = 2$:

Similarly, when $g_0 = 2$, since we assume $0 < e \leq 2d$, $4d - 1 \geq 2d + e - 1 \geq 2^{n-2}d - 2^n$. Then, if $n > 4$,

$$(n \leq) d \leq 4 + \frac{15}{2^{n-2} - 4}. \quad (21)$$

Thus the candidates of (n, d, e, S) are: $(4, 4, 5, 7)$, $(4, 5, 3, 7)$, $(4, 5, 7, 8)$, $(4, 6, 1, 7)$, $(4, 6, 5, 8)$, $(4, 6, 9, 9)$, $(4, 7, 3, 8)$, $(4, 7, 7, 9)$, $(4, 7, 11, 10)$, $(4, 15, 15, 15)$, $(4, 15, 19, 16)$, $(4, 15, 23, 17)$, $(4, 15, 27, 18)$, $(3, 3, 1, 7)$, $(3, 3, 3, 8)$, $(3, 3, 5, 9)$, $(3, 4, 1, 8)$, $(3, 4, 3, 9)$, $(3, 4, 5, 10)$, $(3, 4, 7, 11)$, $(3, 7, 1, 11)$, $(3, 7, 3, 12)$, $(3, 7, 5, 13)$, $(3, 7, 7, 14)$, $(3, 7, 9, 15)$, $(3, 7, 11, 16)$, $(3, 7, 13, 17)$, $(2, 2, 1, 8)$, $(2, 2, 2, 9)$, $(2, 2, 3, 10)$, $(2, 2, 4, 11)$, $(2, 3, 1, 10)$, $(2, 3, 2, 11)$, $(2, 3, 3, 12)$, $(2, 3, 4, 13)$, $(2, 3, 5, 14)$, $(2, 3, 6, 15)$.

• Case $g_0 = 3$:

Next, if $g_0 = 3$ ($0 < e \leq \frac{7}{3}d$), then

$$(5 \leq n \leq) d \leq 4 + \frac{61}{3(2^{n-2} - \frac{16}{3})}. \quad (22)$$

Hence possible (n, d, e, S) are: (5, 8, 17, 9), (4, 4, 9, 9), (4, 5, 6, 9), (4, 5, 10, 10), (4, 6, 3, 9), (4, 6, 7, 10), (4, 6, 11, 11), (4, 7, 4, 10), (4, 7, 8, 11), (4, 7, 12, 12), (4, 7, 16, 13), (4, 15, 4, 16), (4, 15, 8, 17), (4, 15, 12, 18), (4, 15, 16, 19), (4, 15, 20, 20), (4, 15, 24, 21), (4, 15, 28, 22), (4, 15, 32, 23), (3, 3, 2, 9), (3, 3, 4, 10), (3, 3, 6, 11), (3, 4, 1, 10), (3, 4, 3, 11), (3, 4, 5, 12), (3, 4, 7, 13), (3, 4, 9, 14), (3, 7, 2, 15), (3, 7, 4, 16), (3, 7, 6, 17), (3, 7, 8, 18), (3, 7, 10, 19), (3, 7, 12, 20), (3, 7, 14, 21), (3, 7, 16, 22), (2, 2, 1, 10), (2, 2, 2, 11), (2, 2, 3, 12), (2, 2, 4, 13), (2, 3, 1, 13), (2, 3, 2, 14), (2, 3, 3, 15), (2, 3, 4, 16), (2, 3, 5, 17), (2, 3, 6, 18), (2, 3, 7, 19).

4.3.2 Type (B)

• Case $2 \nmid d$:

Now, $d = LCM\{d_i\} \leq \prod d_i \leq \prod(2^{n_i} - 1) < 2^n$. (d_i is the order of $\boxed{\spadesuit}_i$ in (8)). Here, if $g_0 = 1$ ($0 < e \leq d$), then

$$d + e - 1 \leq 2d - 1 < 2^{n+1}. \quad (23)$$

On the other hand, it follows that

$$d + e - 1 \geq 2^{n-2}(q(d) - 4) \quad (24)$$

since $S = 4 + \frac{d+e-1}{2^{n-2}} \geq q(d)$. From (23)(24), one obtains

$$2^{n+1} > 2^{n-2}(q(d) - 4). \quad (25)$$

Consequently, $12 > q(d)$. Besides, we have $20 > q(d)$ for $g_0 = 2$ ($0 < e \leq 2d$) since $2^{n-2}(q(d) - 4) \leq 2d + e - 1 < 2^{n+2}$. By the similar manner, $26 > q(d)$ when $g_0 = 3$ ($0 < e \leq \frac{7}{3}d$).

• Case $2 \mid d$:

In this case, $n_i = l_i m_i, d_i = 2^{r_i} d_i^0$ ($2 \nmid d_i^0$), then $d_i^0 \mid 2^{m_i} - 1$. Let $r := \max\{r_i\}$. Here, we obtain $2^{r_i-1} + 1 \leq l_i \leq 2^{r_i}$ for $r_i \geq 1$. Accordingly, $2^{r-1} + 1 \leq l_1 \leq 2^r$ when we assume l_1 with $r_1 \geq 1$. Now, notice that

$$\boxed{\spadesuit}_i = \begin{pmatrix} \boxed{\star}_i & \boxed{\star}_i & \tilde{O} & \cdots \\ \tilde{O} & \boxed{\star}_i & \ddots & \ddots \\ \vdots & \ddots & \ddots & \boxed{\star}_i \\ \tilde{O} & \cdots & \tilde{O} & \boxed{\star}_i \end{pmatrix} 1 \quad (\boxed{\star}_i) \} m_i. \quad (26)$$

Then

$$d = LCM\{2^{r_i} d_i^0\} = 2^r \cdot LCM\{d_i^0\} \leq 2^r \cdot \prod d_i^0 \quad (27)$$

$$\leq 2^r \cdot \prod (2^{m_i} - 1) \quad (28)$$

$$< \begin{cases} 2^{r+\sum_{i \geq 1} m_i} & (m_1 \geq 2) \\ 2^{r+\sum_{i \geq 2} m_i} & (m_1 = 1). \end{cases} \quad (29)$$

On the other hand, we know

$$dg_0 + e - 1 \geq 2^{n-2}(q(d) - 4). \quad (30)$$

Hence, if $g_0 = 1$ ($0 < e \leq d$), then

$$2d - 1 \geq 2^{n-2}(q(d) - 4). \quad (31)$$

From (29) (31), we obtain

$$2^{r+(\sum_{i \geq 1} m_i)+1} > 2^{n-2}(q(d) - 4) \quad (32)$$

$$2^{3+r+(\sum_{i \geq 1} m_i)-n} > q(d) - 4 \quad (33)$$

$$2^{3+r-2^{r-1}m_1} > q(d) - 4 \quad (34)$$

for $m_1 \geq 2$. Similarly, $2^{3+r-2^{r-1}-1} > q(d) - 4$ for $m_1 = 1$. Therefore, we obtain $8 > q(d)$. In the same way, we have $12 > q(d)$ and $15 > q(d)$ for $g_0 = 2$ and $g_0 = 3$.

From these upper bounds and the property of σ , we obtain a list of possible (g_0, n, d, e, S) .

(1, 4, 6, 3, 6), (2, 5, 12, 9, 8), (2, 5, 12, 17, 9), (2, 5, 14, 13, 9), (2, 5, 14, 21, 10),
(2, 5, 21, 7, 10), (2, 5, 21, 15, 11), (2, 5, 21, 23, 12), (2, 5, 21, 31, 13), (2, 5, 21, 39, 14),
(2, 4, 6, 5, 8), (2, 4, 6, 9, 9), (3, 6, 21, 34, 10), (3, 6, 28, 29, 11), (3, 6, 28, 45, 12),
(3, 6, 28, 61, 13), (3, 5, 21, 2, 12), (3, 5, 21, 10, 13), (3, 5, 21, 18, 14), (3, 5, 21, 26, 15),
(3, 5, 21, 34, 16), (3, 5, 21, 42, 17), (3, 5, 14, 7, 10), (3, 5, 14, 15, 11), (3, 5, 14, 23, 12),
(3, 5, 14, 31, 13), (3, 5, 12, 13, 10), (3, 5, 12, 21, 11), (3, 4, 6, 7, 10), (3, 4, 6, 11, 11).

Next, within the above lists, we construct explicitly classes of hyperelliptic curves C_0/k_d for $g(C_0) \in \{1, 2, 3\}$ such that there is a covering $\pi/k_d : C \rightarrow C_0$ and the covering curve C/k has genus $g(C) = d \cdot g(C_0) + e$ ($e > 0$).

5 Elliptic/Hyperelliptic curves C_0 against GHS attack

5.1 Existence of a model of C over k

Here, we show conditions for existence of a model of C over k .

Consider that C_0 is a hyperelliptic curve over k_d defined by $y^2 = c \cdot f(x)$ where $c \in k_d^\times$, $f(x)$ is a monic polynomial in $k_d[x]$. Denote by $F(x) \in \mathbb{F}_2[x]$ the minimal polynomial of σ . Define $\hat{F}(x) \in \mathbb{F}_2[x]$ as a polynomial such that $x^d + 1 = F(x)\hat{F}(x) \in \mathbb{F}_2[x]$. We have the following necessary and sufficient condition:

C has a model over $k_d \iff$

$$\begin{aligned} F(\sigma)y^2 &\equiv F(\sigma)c = c^{F(q)} \equiv 1 \pmod{(k_d(x)^\times)^2}, \\ G(\sigma)y^2 &\not\equiv 1 \pmod{(k_d(x)^\times)^2} \text{ for } \forall G(x) \mid F(x), G(x) \neq F(x). \end{aligned} \quad (35)$$

Now we know a model of C over k exists iff the extension σ of the Frobenius automorphism $\sigma_{k_d/k}$ is an automorphism of $k_d(C)$ of order d in the separable closure of $k_d(x)$.

Consequently, in the following lemma, we make the condition for c explicitly.

Lemma 5.1. *Assume the condition (35) holds. In order that the curve C has a model over k , c needs to be a square $c \in (k_d^\times)^2$ when $\hat{F}(1) = 0$. When $\hat{F}(1) = 1$, there is a $\phi \in \text{cov}(C/\mathbb{P}^1)$ such that $\sigma\phi$ has order d even if σ does not have order d . Therefore C always has a model over k .*

Proof: Let $M := \left\{ \frac{b(x)}{a(x)} \mid k_d[x] \ni a(x), b(x) : \text{monic} \right\}$.

Now, one has

$$\begin{aligned} F(\sigma)y &\equiv \epsilon c^{\frac{F(q)}{2}} \pmod{M}, & \text{here } \epsilon = \pm 1 \\ \hat{F}(\sigma)F(\sigma)y &\equiv \hat{F}(\sigma)\epsilon c^{\frac{\hat{F}(q)F(q)}{2}} \\ \sigma^{d+1}y &\equiv \epsilon^{\hat{F}(1)} c^{\frac{q^d+1}{2}} \\ \sigma^d y &\equiv \epsilon^{\hat{F}(1)} c^{\frac{q^d-1}{2}} y \end{aligned}$$

We first consider two possibilities of $F(1) = 1$ and $F(1) = 0$ respectively.

- Case $F(1) = 1$:

We notice $\hat{F}(1) = 0$ in this case. From $\sigma^d y \equiv c^{\frac{q^d-1}{2}} y$, it follows that $c^{\frac{q^d-1}{2}} = 1$. Hence $c \in (k_d^\times)^2$.

- Case $F(1) = 0$:

Here, we consider further two possibilities of $\hat{F}(1) = 0$ and $\hat{F}(1) = 1$.

(a) $\hat{F}(1) = 0$

From $\sigma^d y \equiv c^{\frac{q^d-1}{2}} y$, we know $c \in (k_d^\times)^2$.

(b) $\hat{F}(1) = 1$

Then $\sigma^d y \equiv \epsilon c^{\frac{q^d-1}{2}} y$.

If $\epsilon = +1$ and $c \in (k_d^\times)^2$, then σ has order d (i.e. $\sigma^d y = y$).

If $\epsilon = -1$ or $c \notin (k_d^\times)^2$, then σ has order $2d$.

However, we can show that $\exists \phi \in \text{cov}(C/\mathbb{P}^1)$ such that $(\sigma\phi)^d = 1$.
Indeed, suppose $d = 2^r \cdot d_1$ ($2 \nmid d_1$). Since $\sigma\phi = \sigma\phi\sigma^{-1}$, we have

$$\begin{aligned} (\sigma\phi)^d &= \sigma\phi\sigma^{-1} \cdot \sigma^2\phi\sigma^{-2} \cdots \sigma^d\phi\sigma^{-d} \cdot \sigma^d \\ &= \sigma\phi \sigma^2\phi \cdots \sigma^d\phi \sigma^d \\ &= \sigma\phi \sigma^2\phi \cdots \sigma^{2^r d_1}\phi \sigma^d \\ &= (\phi \sigma\phi \sigma^2\phi \cdots \sigma^{2^r-1}\phi)^{d_1} \sigma^d. \end{aligned}$$

Here we use the additive notation of the Galois action on $\text{cov}(C/\mathbb{P}^1) \simeq \mathbb{F}_2^n$. Define

$$J := \left\{ \begin{pmatrix} 0 & 1 & \cdots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ \vdots & & \ddots & 1 \\ 0 & \cdots & 0 & 0 \end{pmatrix} \right\}_{m \leq 2^r}.$$

Then $J^m = O$. Choose $\phi := {}^t(0, 0, \dots, 1)$. Now, $\sigma^i\phi$ corresponds to $(I + J)^i \cdot {}^t(0, \dots, 0, 1)$. Since

$$I + (I + J) + \cdots + (I + J)^{2^r-1} = \begin{cases} O & \text{if } m < 2^r \\ \begin{pmatrix} 0 & \cdots & 0 & 1 \\ 0 & \cdots & 0 & 0 \\ \vdots & \ddots & \vdots & \vdots \\ 0 & \cdots & 0 & 0 \end{pmatrix} & \text{if } m = 2^r, \end{cases}$$

where O is the zero matrix, it follows that

$$\phi \sigma\phi \sigma^2\phi \cdots \sigma^{2^r-1}\phi = \begin{cases} \mathbf{0} = {}^t(0, 0, \dots, 0) & \text{if } m < 2^r \\ \psi := {}^t(1, 0, \dots, 0) & \text{if } m = 2^r. \end{cases}$$

On the other hand, σ^d is an element in the center of $\text{Gal}(k_d(C)/k(x))$, i.e., $\sigma^d \in Z(\text{Gal}(k_d(C)/k(x))) = \{1, \psi\}$. Thus, in the multiplicative notation,

$$(\sigma\phi)^d = (\phi \sigma\phi \sigma^2\phi \cdots \sigma^{2^r-1}\phi)^{d_1} \sigma^d = \begin{cases} 1^{d_1} \cdot 1 = 1 & \text{if } m < 2^r \\ \psi^{d_1} \cdot \psi = 1 & \text{if } m = 2^r. \end{cases}$$

As a result, we can adopt the above $\sigma\phi$ instead of σ .

□

5.2 The defining equations of C_0

Finally, we show how to derive the defining equations of C_0/k_d for candidates of (n, d, g_0, e, S) . Suppose ${}^{F(\sigma)}f(x) \equiv 1 \pmod{(k_d(x)^\times)^2}$ is satisfied. Recall $x^d + 1 = F(x)\hat{F}(x)$. We will define the following notation as $b_i = 1$ when there exists a ramification point $(\alpha^{q^i}, 0)$ on C_0 and let $b_i = 0$ otherwise for $i = 0, \dots, d-1$. Let $\phi(x) := b_{d-1}x^{d-1} + \dots + b_1x + b_0$. We know that $F(x)\phi(x) \equiv 0 \pmod{x^d + 1} \Leftrightarrow \phi(x) \equiv 0 \pmod{\hat{F}(x)}$. Hence $\exists a(x) \in \mathbb{F}_2[x]$, $(a(x), F(x)) = 1$, $\deg a(x) < \deg F(x)$, $\phi(x) \equiv a(x)\hat{F}(x) \pmod{x^d + 1}$ for given n, d .

Further, we define the equivalence $(b_0, b_1, \dots, b_{d-1}) \sim (b_j, \dots, b_{d-1}, b_0, \dots, b_{j-1})$, then corresponding $\phi(x)$'s belong to the same class of C_0 . Indeed, $x^r a(x)\hat{F}(x) \equiv a(x)\hat{F}(x) \pmod{x^d + 1} \Leftrightarrow x^r + 1 \equiv 0 \pmod{\hat{F}(x)}$ for $1 \leq r \leq d$. Since $\hat{F}(x)\mathbb{F}_2[x]/(x^d + 1) \cong \mathbb{F}_2[x]/(F(x))$, the number of the classes of C_0 is $N := \#\{(\mathbb{F}_2[x]/(F(x)))^\times\}/d$.

From the facts, we obtain a procedure to derive the defining equations of C_0 is as follows:

1. Choose a polynomial $a(x) = 1$, then $\phi(x) = \hat{F}(x)$ defines a class of C_0 . If $N = 1$, then this procedure is completed.
2. If $N \neq 1$, choose another polynomial $a(x)$ satisfied the above condition and define $\phi(x) = a(x)\hat{F}(x)$.
3. Find the class of C_0 defined by $\phi(x)$.
4. Repeat step 2,3 until $N - 1$ different polynomials $a(x)$ are found so that the coefficients of $\phi(x)$ defined by $a(x)$ are not cyclic permutation of each others (See the example 5.4 as an instance of $N \neq 1$).

Example 5.1. $n = 2, d = 2$ (Type A)

From $x^2 + 1 = (x + 1)^2$, $F(x) = (x + 1)^2$, $\hat{F}(x) = 1$. Now, choose $a(x) = 1$ since $N = 1$, then $\phi(x) = 1$. Thus, there exists a ramification point $\alpha \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \mid \neq d', 2 \mid d'$) on C_0 .

- Case $g_0 = 2, e = 1, S = 8$

The form of C_0/\mathbb{F}_{q^2} is $y^2 = c \cdot h_2(x)h_1(x)$. Here, $h_1(x) \in \mathbb{F}_q[x]$, $h_2(x) \in \mathbb{F}_{q^2}[x] \setminus \mathbb{F}_q[x]$, $\deg h_2(x) = 2, \deg h_1(x) \in \{4, 3\}$, $c := 1$ or a non-square element in \mathbb{F}_{q^2} because $\hat{F}(1) = 1$. Then $h_2(x) = (x - \alpha_1)(x - \alpha_2)$ since $S = 2/\mathbb{F}_{q^2} + 2/\mathbb{F}_{q^2} + 4/\mathbb{F}_q$ is satisfied (Note that $2/\mathbb{F}_{q^2}$ and $4/\mathbb{F}_q$ mean the numbers of fixed points over \mathbb{F}_{q^2} and \mathbb{F}_q respectively). In this case, notice the ramification points $\alpha_1, \alpha_2 \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q$ or $\alpha_1, \alpha_2 \in \mathbb{F}_{q^4} \setminus \mathbb{F}_{q^2}$, $\alpha_2 := \alpha_1^{q^2}$ (i.e. $d' = 2, 4$). See the list in the Appendix for details.

- Case $g_0 = 1, e = 2, S = 7$

The form of C_0/\mathbb{F}_{q^2} is $y^2 = c \cdot h_2(x)h_1(x)$. Here, $h_1(x) \in \mathbb{F}_q[x]$, $h_2(x) \in \mathbb{F}_{q^2}[x] \setminus \mathbb{F}_q[x]$, $\deg h_2(x) = 3, \deg h_1(x) \in \{1, 0\}$, $c := 1$ or a non-square element in \mathbb{F}_{q^2} . Then $h_2(x) = (x - \alpha_1)(x - \alpha_2)(x - \alpha_3)$ since $S = 2/\mathbb{F}_{q^2} +$

$2/\mathbb{F}_{q^2} + 2/\mathbb{F}_{q^2} + 1/\mathbb{F}_q$. In this case, the ramification points are $\alpha_1, \alpha_2, \alpha_3 \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q$ or $\alpha_1 \in \mathbb{F}_{q^4} \setminus \mathbb{F}_{q^2}$ $\alpha_2 := \alpha_1^{q^2}$ $\alpha_3 \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q$ or $\alpha_1 \in \mathbb{F}_{q^6} \setminus (\mathbb{F}_{q^2} \cup \mathbb{F}_{q^3})$ $\alpha_2 := \alpha_1^{q^2}$ $\alpha_3 := \alpha_1^{q^4}$ (i.e. $d' = 2, 4, 6$).

Example 5.2. $n = 2, d = 3$ (Type A)

$x^3 + 1 = (x + 1)(x^2 + x + 1), F(x) = x^2 + x + 1, \hat{F}(x) = x + 1$. Now, choose $a(x) = 1$ since $N = 1$, then $\phi(x) = x + 1$. Consequently, C_0 has ramification points $\alpha, \alpha^q \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$). However, there is no class of C_0 within the list of the previous section.

Example 5.3. $n = 3, d = 3$ (Type A)

$F(x) = (x + 1)(x^2 + x + 1), \hat{F}(x) = 1$. Similarly, C_0 has a ramification point $\alpha \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$). In this case, consider also $(n, d) = (2, 3)$ (i.e. ramification points $\alpha, \alpha^q \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$)).

• Case $g_0 = 2, e = 3, S = 8$

Then, there exist two cases as follow:

1. $S = 3/\mathbb{F}_{q^3} + 5/\mathbb{F}_q$

C_0/\mathbb{F}_{q^3} is $y^2 = c \cdot (x - \alpha)h_1(x)$. Here, $\alpha \in \mathbb{F}_{q^3}, h_1(x) \in \mathbb{F}_q[x]$, $\deg h_1(x) \in \{5, 4\}$, $c := 1$ or a non-square element in \mathbb{F}_{q^3} .

2. $S = 3/\mathbb{F}_{q^3} + 3/\mathbb{F}_{q^3} + 2/\mathbb{F}_q$

C_0/\mathbb{F}_{q^3} is $y^2 = c \cdot (x - \alpha_1)(x - \alpha_1^q)(x - \alpha_2)(x - \alpha_2^q)h_1(x)$. Here, $(x - \alpha_1)(x - \alpha_1^q)(x - \alpha_2)(x - \alpha_2^q) \in \mathbb{F}_{q^3}[x] \setminus \mathbb{F}_q[x]$, $h_1(x) \in \mathbb{F}_q[x]$, $\deg h_1(x) \in \{2, 1\}$, $c := 1$ or a non-square element in \mathbb{F}_{q^3} . In this case, the ramification points are $\alpha_1, \alpha_2 \in \mathbb{F}_{q^3} \setminus \mathbb{F}_q$ or $\alpha_1 \in \mathbb{F}_{q^6} \setminus (\mathbb{F}_{q^2} \cup \mathbb{F}_{q^3})$ $\alpha_2 := \alpha_1^{q^3}$.

Example 5.4. $n = 4, d = 6$ (Type A)

$x^6 + 1 = (x + 1)^2(x^2 + x + 1)^2$.

• Case(a)

$F(x) = (x^2 + x + 1)^2, \hat{F}(x) = (x + 1)^2$. Now, choose $a(x) = 1$ and $a(x) = x + 1$ since $N = 2$, then $\phi(x) = x^2 + 1$ and $\phi(x) = x^3 + x^2 + x + 1$. In these cases, C_0 has ramification points α, α^{q^2} or $\alpha, \alpha^q, \alpha^{q^2}, \alpha^{q^3} \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 6 \mid d'$).

• Case(b)

$F(x) = (x + 1)^2(x^2 + x + 1), \hat{F}(x) = (x^2 + x + 1)$. Now, choose $a(x) = 1$ since $N = 1$, then $\phi(x) = x^2 + x + 1$. In the case, C_0 has ramification points $\alpha, \alpha^q, \alpha^{q^2} \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 6 \mid d'$).

Example 5.5. $n = 4, d = 6 = 2 \cdot 3$ (Type B)

$x^6 + 1 = (x + 1)^2(x^2 + x + 1)^2, F(x) = (x + 1)^2(x^2 + x + 1)$.

Then we consider the candidates of $(n, d) = (2, 2), (2, 3), (3, 3)$ (i.e. ramification points $\beta \in \mathbb{F}_{q^{e'}} \setminus \mathbb{F}_{q^{l'}}$ ($l' \nmid e', 2 \mid e'$) and $\alpha, \alpha^q \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$) or $\beta \in \mathbb{F}_{q^{e'}} \setminus \mathbb{F}_{q^{l'}}$ ($l' \nmid e', 2 \mid e'$) and $\alpha \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$)).

Lists are shown in the Appendices for all defining equations C_0 .

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Appendices

A Classification for type (A) : $\exists d_i = d$

Here, $h_1(x) \in \mathbb{F}_q[x]$, $h_d(x) \in \mathbb{F}_{q^d}[x] \setminus \mathbb{F}_{q^j}[x]$ ($j \nmid d$),

$\eta := 1$ or a non-square element in \mathbb{F}_{q^d} .

$\alpha, \gamma \in \mathbb{F}_{q^d} \setminus \mathbb{F}_{q^j}$ ($j \nmid d$), $\alpha_i \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', d \mid d'$).

Notice $d' = d, 2d, \dots, \max\{i\}d$ for i in the tables.

$$C_0/k_d : y^2 = c \cdot h(x)h_1(x)$$

$$(1) n = 4, d = 4$$

Then $h(x) = h_d(x)$.

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(4, 4, 1, 1, 5)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})$	1, 0	η
$(4, 4, 2, 5, 7)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})$	3, 2	η
$(4, 4, 3, 9, 9)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})$	5, 4	η
	$(x - \alpha)(x - \gamma^q)(x - \gamma^{q^2})$	5, 4	η

$$(2) n = 4, d = 5$$

$h(x) = h_d(x)$

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(4, 5, 3, 10, 10)$	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)(x - \alpha_i^{q^2})(x - \alpha_i^{q^3})$	0	1

$$(3) n = 4, d = 6$$

$h(x) = h_d(x)$

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(4, 6, 1, 3, 6)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^3})$	0	1
$(4, 6, 2, 9, 9)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})$	3, 2	η

$$(4) n = 4, d = 7$$

$h(x) = h_d(x)$

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(4, 7, 2, 7, 9)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	2, 1	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	2, 1	η
$(4, 7, 2, 11, 10)$	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})$	3, 2	η
	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^3})$	3, 2	η
$(4, 7, 3, 8, 11)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	4, 3	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	4, 3	η
$(4, 7, 3, 12, 12)$	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})$	5, 4	η
	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^3})$	5, 4	η

Here, $\alpha \in \mathbb{F}_{q^d} \setminus \mathbb{F}_{q^j}$ ($j \mid \neq d$), $\alpha_i \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \mid \neq d', d \mid d'$).

$$C_0/k_d : y^2 = c \cdot h(x)h_1(x)$$

$$(5) \ n = 3, d = 3$$

$$h(x) = h_d(x)$$

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
(3, 3, 1, 2, 6)	$x - \alpha$	3, 2	η
(3, 3, 2, 1, 7)	$(x - \alpha)(x - \alpha^q)$	4, 3	η
(3, 3, 2, 3, 8)	$x - \alpha$	5, 4	η
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	2, 1	η
(3, 3, 2, 5, 9)	$(x - \alpha)(x - \alpha^q)(x - \gamma)$	3, 2	η
(3, 3, 3, 2, 9)	$(x - \alpha)(x - \alpha^q)$	6, 5	η
(3, 3, 3, 4, 10)	$x - \alpha$	7, 6	η
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	4, 3	η
(3, 3, 3, 6, 11)	$(x - \alpha)(x - \alpha^q)(x - \gamma)$	5, 4	η
	$\prod_{i=1}^3 (x - \alpha_i)(x - \alpha_i^q)$	2, 1	η

$$(6) \ n = 3, d = 4 \quad \beta \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q, \beta_i \in \mathbb{F}_{q^{e'}} \setminus \mathbb{F}_{q^{l'}} \ (l' \mid \neq e', 2 \mid e'), h_2(x) \in$$

$$\mathbb{F}_{q^2}[x] \setminus \mathbb{F}_q[x]$$

$$h(x) = h_d(x)h_2(x)$$

(n, d, g_0, e, S)	$h_d(x)$	$h_2(x)$	$\deg h_1(x)$	c
(3, 4, 1, 1, 6)	$(x - \alpha)(x - \alpha^q)$	1	2, 1	1
(3, 4, 1, 3, 7)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	1, 0	1
(3, 4, 2, 1, 8)	$(x - \alpha)(x - \alpha^q)$	1	4, 3	1
(3, 4, 2, 3, 9)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	3, 2	1
(3, 4, 2, 5, 10)	$(x - \alpha_1)(x - \alpha_1^q)(x - \alpha_2)(x - \alpha_2^q)$	1	2, 1	1
	$(x - \alpha)(x - \alpha^q)$	$(x - \beta_1)(x - \beta_2)$	2, 1	1
(3, 4, 2, 7, 11)	$(x - \alpha)(x - \alpha^q)$	$\prod_{i=1}^3 (x - \beta_i)$	1, 0	1
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	$x - \beta$	1, 0	1
(3, 4, 3, 1, 10)	$(x - \alpha)(x - \alpha^q)$	1	6, 5	1
(3, 4, 3, 3, 11)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	5, 4	1
(3, 4, 3, 5, 12)	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	1	4, 3	1
	$(x - \alpha)(x - \alpha^q)$	$(x - \beta_1)(x - \beta_2)$	4, 3	1
(3, 4, 3, 7, 13)	$(x - \alpha)(x - \alpha^q)$	$\prod_{i=1}^3 (x - \beta_i)$	3, 2	1
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	$x - \beta$	3, 2	1
(3, 4, 3, 9, 14)	$(x - \alpha)(x - \alpha^q)$	$\prod_{i=1}^4 (x - \beta_i)$	2, 1	1
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	$(x - \beta_1)(x - \beta_2)$	2, 1	1
	$\prod_{i=1}^3 (x - \alpha_i)(x - \alpha_i^q)$	1	2, 1	1

(7) $n = 2, d = 2$
 $h(x) = h_d(x)$

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(2, 2, 1, 1, 6)$	$(x - \alpha_1)(x - \alpha_2)$	2, 1	η
$(2, 2, 1, 2, 7)$	$\prod_{i=1}^3 (x - \alpha_i)$	1, 0	η
$(2, 2, 2, 1, 8)$	$\prod_{i=1}^2 (x - \alpha_i)$	4, 3	η
$(2, 2, 2, 2, 9)$	$\prod_{i=1}^3 (x - \alpha_i)$	3, 2	η

(n, d, g_0, e, S)	$h_d(x)$	$\deg h_1(x)$	c
$(2, 2, 2, 3, 10)$	$\prod_{i=1}^4 (x - \alpha_i)$	2, 1	η
$(2, 2, 2, 4, 11)$	$\prod_{i=1}^5 (x - \alpha_i)$	2, 1	η
$(2, 2, 3, 1, 10)$	$\prod_{i=1}^2 (x - \alpha_i)$	6, 5	η
$(2, 2, 3, 2, 11)$	$\prod_{i=1}^3 (x - \alpha_i)$	5, 4	η
$(2, 2, 3, 3, 12)$	$\prod_{i=1}^4 (x - \alpha_i)$	4, 3	η
$(2, 2, 3, 4, 13)$	$\prod_{i=1}^5 (x - \alpha_i)$	3, 2	η

B Classification for type (B): $\forall d_i \neq d$

Here, $h_1(x) \in \mathbb{F}_q[x]$, $\eta := 1$ or a non-square element in \mathbb{F}_{q^d} .

$C_0/k_d : y^2 = c \cdot h(x)h_1(x)$

(1) $n = 6, d = 28 = 7 \cdot 4$, $\alpha \in \mathbb{F}_{q^7} \setminus \mathbb{F}_q, \beta \in \mathbb{F}_{q^4} \setminus \mathbb{F}_{q^2}$

Then $h(x) = h_7(x)h_4(x)$.

(n, d, g_0, e, S)	$h_7(x)$	$h_4(x)$	$\deg h_1(x)$	c
$(6, 28, 3, 61, 13)$	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	2, 1	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	2, 1	η

(2) $n = 5, d = 12 = 4 \cdot 3$, $\alpha \in \mathbb{F}_{q^4} \setminus \mathbb{F}_{q^2}, \beta \in \mathbb{F}_{q^3} \setminus \mathbb{F}_q$

$h(x) = h_4(x)h_3(x)$

(n, d, g_0, e, S)	$h_4(x)$	$h_3(x)$	$\deg h_1(x)$	c
$(5, 12, 2, 17, 9)$	$(x - \alpha)(x - \alpha^q)$	$(x - \beta)(x - \beta^q)$	2, 1	η
$(5, 12, 3, 21, 11)$	$(x - \alpha)(x - \alpha^q)$	$(x - \beta)(x - \beta^q)$	4, 3	η

(3) $n = 5, d = 14 = 7 \cdot 2$, $\alpha \in \mathbb{F}_{q^7} \setminus \mathbb{F}_q, \beta \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q, \beta_i \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}} (j' \nmid d', 2 \mid d')$, $h_2(x) \in \mathbb{F}_{q^2}[x] \setminus \mathbb{F}_q[x]$

$h(x) = h_7(x)h_2(x)$

(n, d, g_0, e, S)	$h_7(x)$	$h_2(x)$	$\deg h_1(x)$	c
(5, 14, 2, 21, 10)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$x - \beta$	1, 0	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$x - \beta$	1, 0	η
(5, 14, 3, 23, 12)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$x - \beta$	3, 2	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$x - \beta$	3, 2	η
(5, 14, 3, 31, 13)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$\prod_{i=1}^2 (x - \beta_i)$	2, 1	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$\prod_{i=1}^2 (x - \beta_i)$	2, 1	η
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})$	$x - \beta$	4, 3	η
	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^3})$	$x - \beta$	4, 3	η

$$C_0/k_d : y^2 = c \cdot h(x)h_1(x)$$

(4) $n = 5, d = 21 = 7 \cdot 3$, $\alpha \in \mathbb{F}_{q^7} \setminus \mathbb{F}_q, \beta \in \mathbb{F}_{q^3} \setminus \mathbb{F}_q, \beta_i \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$), $h_3(x) \in \mathbb{F}_{q^3}[x] \setminus \mathbb{F}_q[x]$

$$h(x) = h_7(x)h_3(x)$$

(n, d, g_0, e, S)	$h_7(x)$	$h_3(x)$	$\deg h_1(x)$	c
(5, 21, 2, 7, 10)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	0	1
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	0	1
(5, 21, 3, 2, 12)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	2, 1	1
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$(x - \beta)(x - \beta^q)$	2, 1	1
(5, 21, 3, 10, 13)	$(x - \alpha)(x - \alpha^q)(x - \alpha^{q^2})(x - \alpha^{q^4})$	$\prod_{i=1}^2 (x - \beta_i)(x - \beta_i^q)$	0	1
	$(x - \alpha)(x - \alpha^{q^2})(x - \alpha^{q^3})(x - \alpha^{q^4})$	$\prod_{i=1}^2 (x - \beta_i)(x - \beta_i^q)$	0	1

(5) $n = 4, d = 6 = 3 \cdot 2$, $\alpha \in \mathbb{F}_{q^3} \setminus \mathbb{F}_q, \alpha_i \in \mathbb{F}_{q^{d'}} \setminus \mathbb{F}_{q^{j'}}$ ($j' \nmid d', 3 \mid d'$), $\beta \in \mathbb{F}_{q^2} \setminus \mathbb{F}_q, \beta_i \in \mathbb{F}_{q^{e'}} \setminus \mathbb{F}_{q^{l'}}$ ($l' \nmid e', 2 \mid e'$), $h_3(x) \in \mathbb{F}_{q^3}[x] \setminus \mathbb{F}_q[x], h_2(x) \in \mathbb{F}_{q^2}[x] \setminus \mathbb{F}_q[x]$

$$h(x) = h_3(x)h_2(x)$$

(n, d, g_0, e, S)	$h_3(x)$	$h_2(x)$	$\deg h_1(x)$	c
(4, 6, 1, 3, 6)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	0	η
(4, 6, 2, 5, 8)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	3, 2	η
(4, 6, 2, 9, 9)	$x - \alpha$	$x - \beta$	4, 3	η
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	$x - \beta$	1, 0	η
	$(x - \alpha)(x - \alpha^q)$	$\prod_{i=1}^2 (x - \beta_i)$	2, 1	η
(4, 6, 3, 7, 10)	$(x - \alpha)(x - \alpha^q)$	$x - \beta$	5, 4	η
(4, 6, 3, 11, 11)	$x - \alpha$	$x - \beta$	6, 5	η
	$\prod_{i=1}^2 (x - \alpha_i)(x - \alpha_i^q)$	$x - \beta$	3, 2	η
	$(x - \alpha)(x - \alpha^q)$	$\prod_{i=1}^2 (x - \beta_i)$	4, 3	η