Research Letter

Generalized Cumulative Residual Entropy for Distributions with Unrestricted Supports

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We consider the cumulative residual entropy (CRE) a recently introduced measure of entropy. While in previous works distributions with positive support are considered, we generalize the definition of CRE to the case of distributions with general support. We show that several interesting properties of the earlier CRE remain valid and supply further properties and insight to problems such as maximum CRE power moment problems. In addition, we show that this generalized CRE can be used as an alternative to differential entropy to derive information-based optimization criteria for system identification purpose.

1. INTRODUCTION

The concept of entropy is important for studies in many areas of engineering such as thermodynamics, mechanics, or digital communications. An early definition of a measure of the entropy is the Shannon entropy [1, 2]. In Shannon's approach, discrete values and absolutely continuous distributions are treated in a somewhat different way through entropy and differential entropy, respectively. Considering the complementary cumulative distribution function (CCDF) instead of the probability density function in the definition of differential entropy leads to a new entropy measure named cumulative residual entropy (CRE) [3, 4]. In [3, 4], CRE is defined as

$$\mathcal{E}(X) = - \int_{\mathbb{R}^n_+} P(|X| > u) \log P(|X| > u)\,du, \qquad (1)$$

where n is the dimension of the random vector X. Clearly, this formula is valid both for a discrete or an absolutely continuous random variable (RV), or with both a discrete and an absolutely continuous part, because it resorts to the CCDF of $|X|$. In addition, unlike Shannon differential entropy it is always positive, while preserving many interesting properties of Shannon entropy. The concept of CRE has found nice interpretations and applications in the fields of reliability (see

[5] where the concept of dynamic CRE is introduced) and images alignment [3].

Shannon entropy can be seen as a particular case of exponential entropy, when entropy order tends to 1. Thus, following the work in [4], a modified version of the exponential entropy, where PDF is replaced by CCDF, has been introduced in [6], leading to new entropy-type measures, called survival entropies.

However, both Rao et al.'s CRE and its exponential entropy generalization by Zografos and Nadarajah lead to entropy-type definitions that assume either positive valued RVs or apply to $|X|$ otherwise. Although the positive case is of great interest for many applications, CRE and exponential entropies entail difficulties when working with RVs with supports that are not restricted to positive values.

In this paper, we show that for an RV X, (1) remains a valid expression when $P(|X| > u)$ is replaced by $P(X > u)$ and integration is performed over $\mathbb{R}^n$, without further hypothesis than in [4]. In addition, some desirable properties are enabled by this CRE definition extension. We also complete the power moment constrained maximum CRE distributions problem that was adressed in [7], for classes of distributions that have lower-unbounded supports. Finally, we illustrate the potential superiority of the proposed generalized CRE (GCRE) against differential entropy in mutual information-based estimation problems.

The paper is organized as follows. Section 2 introduces the GCRE definition. Some properties of GCRE are discussed in Section 3. In Section 4, we introduce cumulative entropy rate and mutual information rate. Section 5 deals with maximum GCRE distributions. With a view to illustrate the potentiality of GCRE, in Section 6, we show on a simple example a possible benefit of GCRE for systems identification.

2. GENERALIZED CUMULATIVE RESIDUAL ENTROPY (GCRE)

We will denote by $F_X^c(x)$ the complementary cumulative distribution function (survival function) of a multivariate RV $X = [X_1, \ldots, X_n]^T$ of dimension n: $F_X^c(X) = P(X > x) = P(X_i > x_i, \ i = 1, \ldots, n)$. We denote by $H_C(X)$ the GCRE of X that we define by

$$H_C(X) = -\int_{\mathbb{R}^n} F_X^c(u) \log F_X^c(u) du. \tag{2}$$

Clearly, like the CRE, the GCRE is a positive and concave function of F_X^c. In addition, existence of GCRE can be established without further assumption upon distribution than those assumed for the CRE in [4].

Theorem 1. $H_C(X) < \infty$ if for some $p > n$, $E[|X|^p] < \infty$.

Proof. First let us remark that from the proof of the existence of CRE in [4], it is sufficient to prove the result when X is a scalar RV, that is $n = 1$, and for $p > 1$. Then, letting $p^{-1} < \alpha < 1$, we use the following inequality:

$$-F_X^c(x) \log F_X^c(x) \le \frac{e^{-1}}{1-\alpha} [F_X^c(x)]^\alpha 1\mathrm{I}_{]0,\infty[}(x) \tag{3}$$
$$+ (1 - F_X^c(x)) 1\mathrm{I}_{]-\infty,0]}(x),$$

where $1\mathrm{I}_A(x) = 1$ if $x \in A$ and $1\mathrm{I}_A(x) = 0$ otherwise. The existence of $\int_{\mathbb{R}} (e^{-1}/(1-\alpha))[F_X^c(x)]^\alpha 1\mathrm{I}_{[0,\infty[}(x) dx$ can be proven just as in [4]. Now, letting $u = -x$, we have

$$(1 - F_X^c(x)) 1\mathrm{I}_{]-\infty,0]}(x)$$
$$= F_X(x) 1\mathrm{I}_{]-\infty,0]}(x)$$
$$\le 1\mathrm{I}_{[-1,0[}(t) + F_X(x) 1\mathrm{I}_{]-\infty,-1[}(x)$$
$$\le 1\mathrm{I}_{]0,1]}(u) + F_X(-u) 1\mathrm{I}_{]1,\infty]}(u) \tag{4}$$
$$\le 1\mathrm{I}_{]0,1]}(u) + F_{-X}(u) 1\mathrm{I}_{]1,\infty]}(u)$$
$$\le 1\mathrm{I}_{]0,1]}(u) + F_{|-X|}^c(u) 1\mathrm{I}_{]1,\infty]}(u)$$
$$\le 1\mathrm{I}_{]0,1]}(u) + u^{-p} E[|X|^p] 1\mathrm{I}_{]1,\infty]}(u).$$

Thus,

$$\int_{\mathbb{R}} (1 - F_X^c(x)) 1\mathrm{I}_{]-\infty,0]}(x) dx$$
$$\le \int_{\mathbb{R}} (1\mathrm{I}_{]0,1]}(u) + u^{-p} E[|X|^p] 1\mathrm{I}_{]1,\infty]}(u)) du \tag{5}$$
$$\le 1 + \int_1^\infty u^{-p} E[|X|^p] du$$
$$< \infty.$$

Finally, putting all pieces together one finally proves convergence of right-hand side of (2). □

3. A FEW PROPERTIES OF GCRE

Let us now exhibit a few more interesting properties of the GCRE. First, it is easy to check that like Shannon entropy the GCRE remains constant with respect to variable translation:

$$\forall a \in \mathbb{R}^n, \quad H_C(X + a) = H_C(X). \tag{6}$$

In the same way, it is clear that

$$\forall a \in \mathbb{R}_+, \quad H_C(aX) = aH_C(X). \tag{7}$$

When $a < 0$, we do not have such a nice property. However, let us consider the important particular case where the distribution of X has a symmetry of the form

$$\exists \mu, \ \forall x, \quad F_X^c(\mu - x) = 1 - F_X^c(\mu + x). \tag{8}$$

In this case, we get the following result.

Theorem 2. *For an RV X that satisfies symmetry property* (8), *one has*

$$\forall a \in \mathbb{R}, \quad H_C(aX) = |a| H_C(X). \tag{9}$$

Proof. Since it is clear that for all $a \in \mathbb{R}_+$, $H_C(aX) = aH_C(X)$, we just have to check that $H_C(-X) = H_C(X)$, which can be established as follows:

$$-H_C(-X) = \int_{\mathbb{R}} F_{-X}^c(x) \log F_{-X}^c(x) dx$$
$$= \int_{\mathbb{R}} F_{-X}^c(x - \mu) \log F_{-X}^c(x - \mu) dx$$
$$= \int_{\mathbb{R}} F_X(-x + \mu) \log F_X(-x + \mu) dx$$
$$= \int_{\mathbb{R}} (1 - F_X(x + \mu)) \log (1 - F_X(x + \mu)) dx$$
$$= \int_{\mathbb{R}} F_X^c(x + \mu) \log F_X^c(x + \mu) dx$$
$$= \int_{\mathbb{R}} F_X^c(x) \log F_X^c(x) dx$$
$$= -H_C(X). \tag{10}$$
□

When the entries of vector X are independent, it has been shown in [4] that if the X_i are nonnegative, then

$$H_C(X) = \sum_i (\Pi_{j \ne i} \mathbb{E}[X_j]) H_C(X_i). \tag{11}$$

However, this formula does not extend to RVs with distributions carried by $\mathbb{R}^n$ because $F^C(X)$ can be integrated over $\mathbb{R}_+^n$ in general but never over $\mathbb{R}^n$. However, if the X_is

are independent and have lower bounded supports with respective lower bounds $m_1, \ldots, m_n$,

$$
\begin{aligned}
H_C(X) &= \int_{\Pi_i]m_i, \infty]} -F_X^c(x) \log F_X^c(x) dx \\
&= \sum_i (\Pi_{j \neq i}(\mathbb{E}[X_j] - m_j)) H_C(X_i),
\end{aligned}
\tag{12}
$$

because

$$
\begin{aligned}
\int_{m_i}^{\infty} F_{X_i}^c(u) du &= [u F_{X_i}^c(u)]_{m_i}^{\infty} + \int_{m_i}^{\infty} u P_{X_i}(du) \\
&= -m_i + \mathbb{E}[X_i].
\end{aligned}
\tag{13}
$$

Conditional GCRE definition is a direct extension of the definition of conditional CRE: the conditional GCRE of X knowing that Y is equal to y is defined by

$$
\begin{aligned}
&H_C(X | Y = y) \\
&= \int_{\mathbb{R}^n} -P(X > x | Y = y) \log P(X > x | Y = y) dx.
\end{aligned}
\tag{14}
$$

We recall an important result from [4] that states that conditioning reduces the entropy.

Theorem 3. *For any X and Y,*

$$
H_C(X | Y) \le H_C(X)
\tag{15}
$$

equality holds if and only if X is independent of Y.

As a consequence, if $X \to Y \to Z$ is a Markov chain, we have the data processing inequality for GCRE:

$$
H_C(Z | X, Y) \le H_C(Z | X).
\tag{16}
$$

4. ENTROPY AND MUTUAL INFORMATION RATES

4.1. Entropy rate

The GCRE of a stochastic process $\{X_i\}$ is defined by

$$
H_C(X) = \lim_{n \to \infty} H_C(X_n | X_{n-1}, X_{n-2}, \ldots, X_1)
\tag{17}
$$

when the limit exists.

Theorem 4. *For stationary processes, the limit exists.*

Proof. Consider

$$
\begin{aligned}
H_C(X_{n+1} | X_n, \ldots, X_1) &\le H_C(X_{n+1} | X_n, \ldots, X_2) \\
&\le H_C(X_n | X_{n-1}, \ldots, X_1).
\end{aligned}
\tag{18}
$$

The first line follows from the fact that conditioning reduces entropy and the second follows from the stationarity (see [2] for the equivalent proof in the case of Shannon entropy). □

4.2. Mutual information

Let X and Y be two RVs. We define the cumulative mutual information between X and Y as follows:

$$
I_C(X; Y) = H_C(X) - H_C(X | Y).
\tag{19}
$$

Theorem 5. *I_C is nonnegative and it vanishes if and only if X and Y are independent.*

Proof. It is clear that I_C is nonnegative because of Theorem 3. □

For a random vector $\mathbf{X} = (X_1, X_2, \ldots, X_n)$ of size n, mutual information is defined by

$$
I_C(\mathbf{X}) = \sum_{i=1}^{n} H_C(X_i) - H_C(X_n | X_{n-1}, \ldots, X_1).
\tag{20}
$$

In the case of stochastic processes $\{X_i\}$, we have $H_C(X) = \lim_{n \to \infty} H_C(X_n | X_{n-1}, \ldots, X_1)$ and the limit exists for stationary processes. Then the mutual information rate for $\{X_i\}$ is defined as

$$
I_C(X) = \lim_{T \to \infty} \frac{1}{T} \sum_{t=1}^{T} H_C(X_t) - H_C(X),
\tag{21}
$$

where $H_C(X_t)$ is the marginal GCRE of the process X.

5. MAXIMUM GCRE DISTRIBUTIONS

In this section, we only consider the case of one-dimensional RVs ($n = 1$). Maximum entropy principle is useful in many scientific areas and most important distributions can be derived from it [8]. The maximum CRE distribution has been studied in [7]. For an RV X with a symmetric CCDF in the sense of (8), we are looking for the maximum GCRE distribution, that is, the CCDF that solves the following moment problem:

$$
\begin{aligned}
&\max_{F^c} H_C(F^c) \\
&\int_{\mathbb{R}} r_i(x) p(x) dx = c_i, \quad i = 1, \ldots, m,
\end{aligned}
\tag{22}
$$

where $p(x) = -(d/dx)F^c(x)$, $(r_i)_{i=1,m}$, and $(c_i)_{i=1,m}$ are fixed C^1 real valued functions and real coefficients, respectively. The solution of this problem is supplied by the following result.

Theorem 6. *When the symmetry property (8) holds, the solution of problem (22), when it can be reached, is of the form*

$$
F^c(x) = \frac{1}{1 + \exp [\sum_1^m \lambda_i r_i'(x - \mu)]} \quad (x \ge \mu).
\tag{23}
$$

Proof.

$$-H_C(X) = \int_{-\infty}^{\mu} F^c(x) \log (F^c(x)) dx$$

$$+ \int_{\mu}^{\infty} F^c(x) \log (F^c(x)) dx$$

$$= \int_{\mathbb{R}_-} F^c(x + \mu) \log (F^c(x + \mu)) dx$$

$$+ \int_{\mathbb{R}_+} F^c(x + \mu) \log (F^c(x + \mu)) dx \qquad (24)$$

$$= \int_{\mathbb{R}_-} (1 - F^c(\mu - x)) \log (1 - F^c(\mu - x)) dx$$

$$+ \int_{\mathbb{R}_+} F^c(x + \mu) \log (F^c(x + \mu)) dx$$

$$= \int_{\mathbb{R}_+} \big[(1 - F^c(\mu + x)) \log (1 - F^c(\mu + x))$$

$$+ F^c(\mu + x) \log (F^c(\mu + x)) \big] dx.$$

Let us define $\tilde{f}$ by

$$\tilde{f}(F^c(x), -p)$$

$$= -(1 - F^c(\mu + x)) \log (1 - F^c(\mu + x))$$

$$- F^c(\mu + x) \log (F^c(\mu + x)) + \sum_{1}^{m} \lambda_i (-p(x)) r_i(x). \qquad (25)$$

Then, since $(F^c(x))' = -p(x)$, the Euler-Lagrange equation [9] states that the solution F^c of problem (22) is a solution of equation

$$\frac{d}{dx} \tilde{f}_{-p}(x) = \tilde{f}_{F^c}(x), \qquad (26)$$

where $\tilde{f}_u$ is the partial derivative of $\tilde{f}$ with respect to component u. From (25), we get

$$\frac{d}{dx} \tilde{f}_{-p}(x) = \sum_{1}^{m} \lambda_i r_i'(x),$$

$$\tilde{f}_{F^c}(x) = \log (1 - F^c(\mu + x)) - \log F^c(\mu + x). \qquad (27)$$

Then,

$$\log \frac{1 - F^c(\mu + x)}{F^c(\mu + x)} = \sum_{1}^{m} \lambda_i r_i'(x),$$

$$F^c(x) = \frac{1}{1 + \exp \left(\sum_{1}^{m} \lambda_i r_i'(x - \mu) \right)}, \qquad (28)$$

for $x \in [\mu, \infty[$. $\qquad \square$

5.1. Example

We set the constraints $\mathbb{E}[X] = \mu$ and $\mathbb{E}[X^2] = \sigma^2$. Then the maximum GCRE symmetric solution for the CCDF of X is given by

$$F_X^c(x) = \frac{1}{1 + \exp (\lambda_1 + 2\lambda_2 x)}, \qquad (29)$$

for $x > 0$, which is the CCDF of a logistic distribution. The moment constraints lead to $\lambda_1 + 2\lambda_2 x = (\sigma\sqrt{3}/\pi)(x - \mu)$. The corresponding PDF is defined on $\mathbb{R}$ by

$$p_X(x) = \frac{\sigma\sqrt{3}}{\pi} \frac{\exp \left(-(\sigma\sqrt{3}/\pi)(x - \mu) \right)}{\left(1 + \exp \left(-(\sigma\sqrt{3}/\pi)(x - \mu) \right) \right)^2}. \qquad (30)$$

5.2. Positive random variables

It has been shown in [7] that the maximum CRE (i.e., the maximum GCRE under additional nonnegative constraint) distribution has CCDF in the form

$$F^C(x) = \exp \left(-\sum_{i=1}^{m} \lambda_i r_i'(x) \right), \qquad (31)$$

for $x \in [0, \infty[$. In [7], this result is derived from the log-sum inequality, but of course it can also be derived from the Euler-Lagrange equation along the same lines as in the proof of Theorem 6.

With a positive support constraint and under first and second moment constraints, it comes that the optimum CCDF is of the form $F^c(x) = \exp(-\lambda_1 - 2\lambda_2 x)$, for $x > 0$. Thus the solution, if it exists, is an exponential distribution. In fact, the first and second power moment constraints must be such that $\mathbb{E}[X^2] = 2(\mathbb{E}[X])^2$, otherwise the problem has no exact solution.

6. SIMULATION RESULTS

With a view to emphasize the potential practical interest of GCRE, we consider a simple system identification problem. Here, we consider an $MA(1)$ process, denoted by $Y = (Y_n)_{n \in \mathbb{Z}}$, generated by a white noise $X = (X_n)_{n \in \mathbb{Z}}$ and corrupted by a white noise N:

$$Y_n = X_n - aX_{n-1} + N_n. \qquad (32)$$

The model input X and output Y are observed and the system model $(MA(1))$ is assumed to be known. We want to estimate the coefficient a without prior knowledge upon the distributions of X and Y. Thus, we resort to mutual information (MI) to estimate a as the coefficient α such that RVs $\hat{Y}_n^\alpha = X_n - \alpha X_{n-1}$ and Y_n show the highest dependence. Shannon MI between $\hat{Y}_n^\alpha$ and Y_n is given by $f_S(\alpha) = I_S(\hat{Y}_n^\alpha, Y_n) = H_S(\hat{Y}_n^\alpha) - H_S(\hat{Y}_n|Y_n)$, where H_S is Shannon differential entropy. Similarly, for GCRE, MI will be defined as $f_C(\alpha) = I_C(\hat{Y}_n^\alpha, Y) = H_C(\hat{Y}_n^\alpha) - H_C(\hat{Y}_n^\alpha|Y_n)$. We compare estimation performance for a by maximizing both $f_S(\alpha)$ and $f_C(\alpha)$. Since true values of $f_S(\alpha)$ and $f_C(\alpha)$ are not available, they are estimated from empirical distributions of $(Y_n, \hat{Y}_n^\alpha)$.

For simulations, we have chosen X Gaussian and N with a Laplace distribution: $p_N(x) = (\lambda/2) \exp(-\lambda|x|)$. We consider an experiment with $a = 0.5$ and noise variance equal to 0.2. Estimation is carried out from observation of $(X_n, Y_n)_{n=1,400}$. Here, optimization of MIs is realized on a fixed regular grid of 200 points over interval [0,1]. Estimation performance is calculated from 200 successive
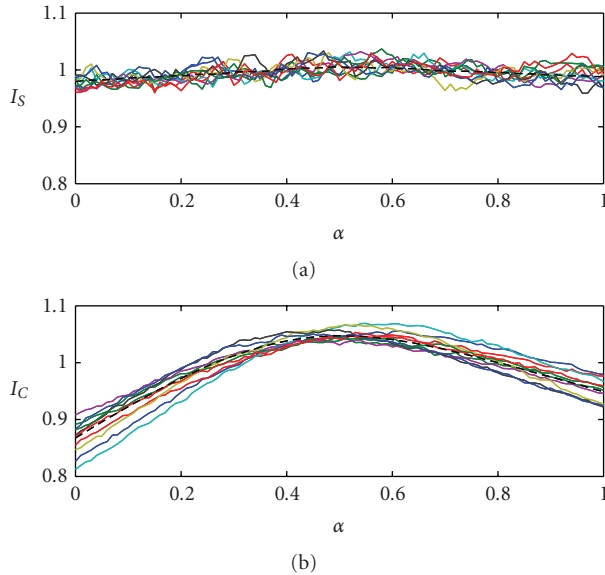
I_S

(a)

I_C

(b)

FIGURE 1: Ten estimates of (a) $f_S(\alpha)$ and (b) $f_C(\alpha)$. Dotted line: mean estimate averaged from 200 realizations.

experiments. Estimation of a from Shannon MI leads to bias and standard deviation that are equal to 0.032 and 0.18, respectively, while they are equal to 0.004 and 0.06, respectively, for GCRE MI.

More important, we see on Figure 1(a) that Shannon MI estimates are much more irregular than GCRE MI (Figure 1(b)) estimates because of smoothing brought by density integration in the calculation of CCDF. This difference is important since the use of an iterative local optimization technique would have failed in general to find Shannon's estimated MI global optimum, because of its many local maxima.

Of course, this drawback can be partly solved by kernel smoothing of the empirical distribution of $(Y_n, \hat{Y}_n^\alpha)$, for instance by using the method proposed in [10]. However, we have checked that, for the above example, very strong smoothing is necessary and then bias and variance performance remain worse than with GCRE MI estimator.

7. CONCLUSION

We have shown that the concept of cumulative residual entropy (CRE) introduced in [3, 4] can be extended to distributions with general supports. Generalized CRE (GCRE) shares many nice features of CRE. We also pointed out specific properties of GCRE such as its maximum, moment constrained, and distribution and we have illustrated practical interest of GCRE by showing how it can be used in system identification procedures.

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The main idea in network coding was introduced in 2000 by Ahlswede et al. With network coding, an intermediate node in a network cannot only forward its incoming packets but also encode them. It has been shown that the use of network coding can enhance the performance of wired networks significantly. Recent work indicates that network coding can also offer significant benefits for wireless networks.

Communications over wireless channels are error-prone and unpredictable due to fading, mobility, and intermittent connectivity. Moreover, in wireless networks, transmissions are broadcast and can be overhead by neighbors, which is treated in current systems as interference. Finally, security poses new challenges in wireless networks, where both passive and active attacks have quite different premises than in wired networks. Ideas in network coding promise to help toward all these issues, allowing to gracefully add redundancy to combat errors, take advantage of the broadcast nature of the wireless medium and achieve opportunistic diversity, exploit interference rather than be limited by it, and provide secure network communication against adversarial attacks.

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Advanced Image Processing for Defense and Security Applications

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The history of digital image processing can be traced back to the 1920s when digital images were transferred between London and New York. However, in the past, the cost of processing was very high because the imaging sensors and computational equipments were very expensive and had only limited functions. As a result, the development of digital image processing was limited.

As optics, imaging sensors, and computational technology advanced, image processing has become more commonly used in many different areas. Some areas of application of digital image processing include image enhancement for better human perception, image compression and transmission, as well as image representation for automatic machine perception.

Most notably, digital image processing has been widely deployed for defense and security applications such as small target detection and tracking, missile guidance, vehicle navigation, wide area surveillance, and automatic/aided target recognition. One goal for an image processing approach in defense and security applications is to reduce the workload of human analysts in order to cope with the ever increasing volume of image data that is being collected. A second, more challenging goal for image processing researchers is to develop algorithms and approaches that will significantly aid the development of fully autonomous systems capable of decisions and actions based on all sensor inputs.

The aim of this special issue is to bring together researchers designing or developing advanced image processing techniques/systems, with a particular emphasis on defense and security applications. Prospective papers should be unpublished and present innovative research work offering contributions either from a methodological or application point of view. Topics of interest include, but are not limited to:

- Multispectral/hyperspectral image processing for object tracking and classification with emphasis on defense-related targets and objects
- Real-time image processing for surveillance, reconnaissance, and homeland security
- Biometric image processing for personal authentication and identification with emphasis on homeland security applications

- Image encryption for secure image storage and transmission
- Image processing to enable autonomous and intelligent control for military, intelligence, and homeland security applications
- Image processing for mental workload evaluation with emphasis on homeland security applications
- Image interpolation and registration for object visualization, tracking, and/or classification

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