

ON THE COMPUTATION OF EDIT DISTANCE FUNCTIONS

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ABSTRACT. The edit distance between two graphs on the same labeled vertex set is the symmetric difference of the edge sets. The edit distance function of hereditary property, \mathcal{H} , is a function of $p \in [0, 1]$ and is the limit of the maximum normalized distance between a graph of density p and \mathcal{H} .

This paper uses localization, for computing the edit distance function of various hereditary properties. For any graph H , $\text{Forb}(H)$ denotes the property of not having an induced copy of H . We compute the edit distance function for $\text{Forb}(H)$, where H is any so-called split graph, and the graph H_9 , a graph first used to describe the difficulties in computing the edit distance function.

1. INTRODUCTION

This paper uses the method of localization, introduced in [12] as a way to compute edit distance functions. It uses some properties of quadratic programming, first applied by Marchant and Thomason [11]. Some results on the edit distance function can be found in a variety of papers [15, 5, 6, 1, 2, 3, 4, 10, 11, 13, 14]. Much of the background to this paper can be found in a paper by Balogh and the author. Terminology and proofs of supporting lemmas that are suppressed here can be found in [12].

1.1. The edit distance function. A **hereditary property** is a family of graphs that is closed under isomorphism and the taking of induced subgraphs. The **edit distance function** of a hereditary property \mathcal{H} , denoted $ed_{\mathcal{H}}(p)$, measures the maximum distance of a density p graph from a hereditary property. Formally, if $\text{Dist}(G, \mathcal{H}) = \min\{|E(G) \Delta E(G')| : |V(G')| = n, G' \in \mathcal{H}\}$, then

$$(1) \quad ed_{\mathcal{H}}(p) = \lim_{n \rightarrow \infty} \max \left\{ \text{Dist}(G, \mathcal{H}) : |V(G)| = n, |E(G)| = \left\lfloor p \binom{n}{2} \right\rfloor \right\} / \binom{n}{2}.$$

In [7], a result of Alon and Stav [1] is generalized to show that the limit in (1) does indeed exist for nontrivial hereditary properties and, furthermore,

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that

$$ed_{\mathcal{H}}(p) = \lim_{n \rightarrow \infty} \text{Dist}(G(n, p), \mathcal{H}) / \binom{n}{2}.$$

For any nontrivial hereditary property \mathcal{H} (that is, one that is not finite), the function $ed_{\mathcal{H}}(p)$ is continuous and concave down. Hence, it achieves its maximum at a point $(p_{\mathcal{H}}^*, d_{\mathcal{H}}^*)$. It should be noted that, for some hereditary properties, $p_{\mathcal{H}}^*$ might be an interval.

1.2. Main results. The main results of this paper are Theorem 1 and Theorem 3.

A **split graph** is a graph whose vertex set can be partitioned into one clique and one independent set. If H is a split graph on h vertices with independence number α and clique number ω , then $\alpha + \omega \in \{h, h + 1\}$. The value of (p^*, d^*) had been obtained for the claw by Alon and Stav [2] and for graphs of the form $K_a + E_b$ (an a -clique with b isolated vertices) by Balogh and the author [7].

Theorem 1. *Let H be a split graph that is neither complete nor empty, with independence number α and clique number ω . Then,*

$$(2) \quad ed_{\text{Forb}(H)}(p) = \min \left\{ \frac{p}{\omega - 1}, \frac{1 - p}{\alpha - 1} \right\}.$$

It is a trivial result (see, e.g., [12]) that $ed_{\text{Forb}(K_{\omega})}(p) = p/(\omega - 1)$ and $ed_{\text{Forb}(\overline{K}_{\alpha})}(p) = (1 - p)/(\alpha - 1)$. So, we can combine Theorem 1 with the prior results for which H is either complete or empty.

Corollary 2. *Let H be a split graph with independence number α and clique number ω . Then, $(p_{\mathcal{H}}^*, d_{\mathcal{H}}^*) = \left(\frac{\omega - 1}{\alpha + \omega - 2}, \frac{1}{\alpha + \omega - 2} \right)$.*

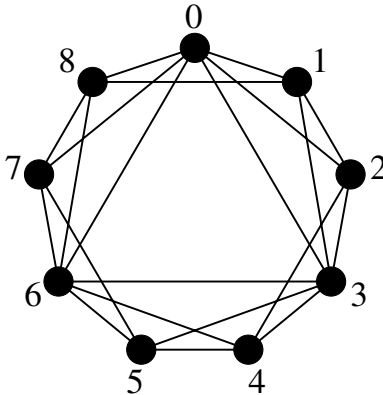


FIGURE 1. The graph H_9 .

The graph, H_9 , as drawn in Figure 1.2, was given in [7] as an example of a hereditary property $\mathcal{H} = \text{Forb}(H_9)$ such that the maximum value of $ed_{\mathcal{H}}(p)$ cannot be determined by CRGs that only have gray edges. In [7] only an

upper bound of $\min \left\{ \frac{p}{3}, \frac{p}{2+2p}, \frac{1-p}{2} \right\}$ is provided for $ed_{\text{Forb}(H_9)}(p)$. Here we determine the function itself.

Theorem 3. *Let H_9 be the graph in Figure 1.2. Then,*

$$ed_{\text{Forb}(H_9)}(p) = \min \left\{ \frac{p}{3}, \frac{p}{1+4p}, \frac{1-p}{2} \right\}.$$

Consequently, $\left(p_{\text{Forb}(H_9)}^*, d_{\text{Forb}(H_9)}^* \right) = \left(\frac{1+\sqrt{17}}{8}, \frac{7-\sqrt{17}}{16} \right)$.

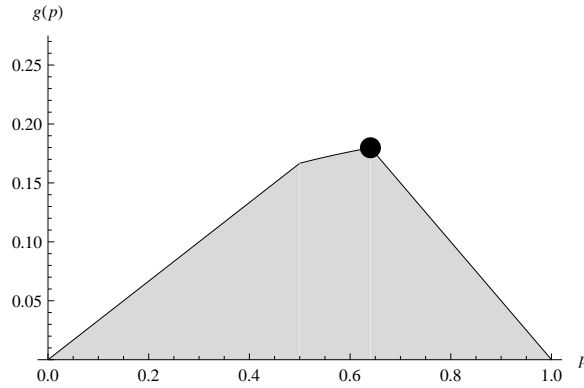


FIGURE 2. Plot of $ed_{\text{Forb}(H_9)}(p) = \min\{p/3, p/(1+4p), (1-p)/2\}$. The point $(p^*, d^*) = \left(\frac{1+\sqrt{17}}{8}, \frac{7-\sqrt{17}}{16} \right)$ is indicated.

The rest of the paper is organized as follows: Section 2 gives some of the general definitions for the edit distance function, such as colored regularity graphs. Section 3 defines and categorizes so-called p -core colored regularity graphs introduced by Marchant and Thomason [11]. Section 5 proves Theorem 1 regarding split graphs. Section 6 proves Theorem 3 regarding the graph H_9 . Section 7 is a section of acknowledgements.

2. BACKGROUND AND BASIC FACTS

2.1. Notation. All graphs are simple. If S and T are sets, then $S + T$ denotes the disjoint union of S and T . If v and w are adjacent vertices in a graph, we denote the edge between them to be vw .

2.2. Colored regularity graphs. A **colored regularity graph (CRG)**, K , is a simple complete graph, together with a partition of the vertices into black and white $V(K) = \text{VW}(K) + \text{VB}(K)$ and a partition of the edges into black, white and gray $E(K) = \text{EW}(K) + \text{EG}(K) + \text{EB}(K)$. We say that a graph H embeds in K , (writing $H \mapsto K$) if there is a function $\varphi : V(H) \rightarrow V(K)$ so that if $h_1 h_2 \in E(H)$, then either $\varphi(h_1) = \varphi(h_2) \in \text{VB}(K)$ or $\varphi(h_1)\varphi(h_2) \in \text{EB}(K) \cup \text{EG}(K)$ and if $h_1 h_2 \notin E(H)$, then either $\varphi(h_1) = \varphi(h_2) \in \text{VW}(K)$ or $\varphi(h_1)\varphi(h_2) \in \text{EW}(K) \cup \text{EG}(K)$.

For a hereditary property of graphs, we denote $\mathcal{K}(\mathcal{H})$ to be the subset of CRGs such that no forbidden graph maps into K . That is, if $\mathcal{F}(\mathcal{H})$ is defined so that $\mathcal{H} = \bigcap_{H \in \mathcal{F}(\mathcal{H})} \text{Forb}(H)$, then $\mathcal{K}(\mathcal{H}) = \{K : H \not\rightarrow K, \forall H \in \mathcal{F}(\mathcal{H})\}$. A CRG K' is said to be a **sub-CRG of K** if K' can be obtained by deleting vertices of K .

2.3. The f and g functions. For every CRG, K , we associate two functions. The function f is a linear function of p and g is found by weighting the vertices. Let K have a total of k vertices $\{v_1, \dots, v_k\}$, and let $\mathbf{M}_K(p)$ be a matrix such that the entries are:

$$[\mathbf{M}_K(p)]_{ij} = \begin{cases} p, & \text{if } v_i v_j \in \text{VW}(K) \cup \text{EW}(K); \\ 1 - p, & \text{if } v_i v_j \in \text{VB}(K) \cup \text{EB}(K); \\ 0, & \text{if } v_i v_j \in \text{EG}(K). \end{cases}$$

Then, we can express the f and g functions over the domain $p \in [0, 1]$ as follows, with $\text{VW} = \text{VW}(K)$, $\text{VB} = \text{VB}(K)$, $\text{EW} = \text{EW}(K)$ and $\text{EB} = \text{EB}(K)$:

$$(3) \quad f_K(p) = \frac{1}{k^2} [p(|\text{VW}| + 2|\text{EW}|) + (1-p)(|\text{VB}| + 2|\text{EB}|)]$$

$$(4) \quad g_K(p) = \begin{cases} \min \mathbf{x}^T \mathbf{M}_K(p) \mathbf{x} \\ \text{s.t. } \mathbf{x}^T \mathbf{1} = 1 \\ \mathbf{x} \geq \mathbf{0} \end{cases}$$

If we denote $\mathbf{1}$ to be the vector of all ones, then $f_K(p) = \left(\frac{1}{k}\mathbf{1}\right)^T \mathbf{M}_K(p) \left(\frac{1}{k}\mathbf{1}\right)$. So, $f_K(p) \geq g_K(p)$.

Theorem 4 ([7]). *For any nontrivial hereditary property \mathcal{H} ,*

$$ed_{\mathcal{H}}(p) = \lim_{K \in \mathcal{K}(\mathcal{H})} g_K(p) = \lim_{K \in \mathcal{K}(\mathcal{H})} f_K(p).$$

2.4. Basic observations on $ed_{\mathcal{H}}(p)$. The following is a summary of basic facts about the edit distance function. Item (iii) comes from Alon and Stav [1]. Item (iv) comes from [7].

Theorem 5. *Let \mathcal{H} be a nontrivial hereditary property with chromatic number χ , complementary chromatic number $\bar{\chi}$, binary chromatic number χ_B and edit distance function $ed_{\mathcal{H}}(p)$.*

- (i) *If $\chi > 1$, then $ed_{\mathcal{H}}(p) \leq p/(\chi - 1)$.*
- (ii) *If $\bar{\chi} > 1$, then $ed_{\mathcal{H}}(p) \leq (1-p)/(\bar{\chi} - 1)$.*
- (iii) *$ed_{\mathcal{H}}(1/2) = 1/(2(\chi_B - 1))$.*
- (iv) *$ed_{\mathcal{H}}(p)$ is continuous and concave down.*
- (v) *$ed_{\mathcal{H}}(p) = ed_{\bar{\mathcal{H}}}(1-p)$.*

3. THE p -CORES

In Marchant and Thomason [11], it is shown that

$$ed_{\mathcal{H}}(p) = \inf \{g_K(p) : K \in \mathcal{K}(\mathcal{H})\} = \inf \{f_K(p) : K \in \mathcal{K}(\mathcal{H})\}.$$

Although the setting of that paper is not edit distance, the results can be translated to our setting. They show, in fact, that $ed_{\mathcal{H}}(p) = \min \{g_K(p) : K \in \mathcal{K}(\mathcal{H})\}$. That is, for any hereditary property \mathcal{H} and $p \in [0, 1]$, there is a CRG, $K \in \mathcal{K}(\mathcal{H})$ such that $ed_{\mathcal{H}}(p) = g_K(p)$. This is found by looking at so-called p -cores. A CRG, K , is a p -core CRG, or simply a p -core, if $g_K(p) < g_{K'}(p)$ for all nontrivial sub-CRGs K' of K . Marchant and Thomason prove that

$$ed_{\mathcal{H}}(p) = \min \{g_K(p) : K \in \mathcal{K}(\mathcal{H}) \text{ and } K \text{ is } p\text{-core}\}.$$

4. COMPUTING EDIT DISTANCE FUNCTIONS USING LOCALIZATION

Upper bounds for the edit distance function of \mathcal{H} are found by simply exhibiting some CRGs $K \in \mathcal{K}(\mathcal{H})$ and computing $g_K(p)$ by means of (4). The localization method obtains lower bounds for $ed_{\mathcal{H}}(p)$. We have already seen much of the theoretical underpinnings. We combine the observations below:

Lemma 6. *Let \mathcal{H} be a nontrivial hereditary property and $p \in (0, 1)$, $\mathcal{K}(\mathcal{H})$ the set of CRGs defined by \mathcal{H} and $\mathcal{K}_p(\mathcal{H})$ the set of p -core CRGs defined by \mathcal{H} . Then,*

- (i) $ed_{\mathcal{H}}(p) = \min\{g_K(p) : K \in \mathcal{K}(\mathcal{H}) \text{ and } K \text{ is } p\text{-core}\}$.
- (ii) *If $p \leq 1/2$ and K is a p -core CRG, then K has no black edges and white edges can only be incident to black vertices.*
- (iii) *If $p \geq 1/2$ and K is a p -core CRG, then K has no white edges and black edges can only be incident to white vertices.*
- (iv) *If \mathbf{x} is the optimal weight function of a p -core CRG K , then for all $v \in V(K)$, $g_K(p) = pd_W(v) + (1-p)d_B(v)$.*

The overall idea is that we need only consider p -core CRGs and their special structure, then a great deal of information can be obtained by focusing on a single vertex. This is referred to as ‘‘localization’’ because we can focus on one vertex at a time.

Lemma 7 has all of the elements to express $d_G(v)$ for any vertex v in a p -core CRG. It is often useful to focus on the gray neighborhood of vertices.

Lemma 7 (Localization). *Let $p \in (0, 1)$ and K be a p -core CRG with optimal weight function \mathbf{x} .*

- (i) *If $p \leq 1/2$, then, $\mathbf{x}(v) = g_K(p)/p$ for all $v \in VW(K)$ and*

$$d_G(v) = \frac{p - g_K(p)}{p} + \frac{1 - 2p}{p}\mathbf{x}(v), \quad \text{for all } v \in VB(K).$$

- (ii) *If $p \geq 1/2$, then $\mathbf{x}(v) = g_K(p)/(1-p)$ for all $v \in VB(K)$ and*

$$d_G(v) = \frac{1 - p - g_K(p)}{1 - p} + \frac{2p - 1}{1 - p}\mathbf{x}(v), \quad \text{for all } v \in VW(K).$$

Corollary 8. *Let $p \in (0, 1)$ and K be a p -core CRG with optimal weight function \mathbf{x} .*

- (i) If $p \leq 1/2$, then $\mathbf{x}(v) \leq g_K(p)/(1-p)$ for all $v \in \text{VB}(K)$.
- (ii) If $p \geq 1/2$, then $\mathbf{x}(v) \leq g_K(p)/p$ for all $v \in \text{VW}(K)$.

Remark 9. From this point forward in the paper, if K is a CRG under consideration and p is fixed, $\mathbf{x}(v)$ will denote the weight of $v \in V(K)$ under the optimal solution of the quadratic program in equation (4) that defines g_K .

One more useful observation is Theorem 6 from [12]:

Theorem 10. A sub-CRG, K' , of a CRG, K , is a **component** if, for all $v \in V(K')$ and all $w \in V(K) - V(K')$, then vw is gray. Let K be a CRG with components $K^{(1)}, \dots, K^{(\ell)}$. Then

$$(g_K(p))^{-1} = \sum_{i=1}^{\ell} (g_{K^{(i)}}(p))^{-1}.$$

5. Forb(H), H A SPLIT GRAPH

We need to define a special class of graphs. For $\omega \geq 2$ and a nonnegative integer vector $(\omega; a_0, a_1, \dots, a_\omega)$, a $(\omega; a_0, a_1, \dots, a_\omega)$ -**clique-star**¹ is a graph G such that $V(G)$ is partitioned into A and W . The set A induces an independent set, the set $W = \{w_1, \dots, w_\omega\}$ induces a clique and for $i = 1, \dots, \omega$, vertex w_i is adjacent to a set of $a_i + 1$ leaves in A and there are a_0 independent vertices. Note that this implies that $\sum_{j=0}^{\omega} a_j = \alpha - \omega$.

Colloquially, a clique-star can be partitioned into stars and independent sets such that the centers of the stars are connected by a clique and there are no other edges. (If one of the stars is K_2 , one of the endvertices is designated to be the center.) Proving that Theorem 1 is true is much more difficult in the case where either H or its complement is a clique-star.

5.1. Proof of Theorem 1. Note that, because H is neither complete nor empty, $\alpha, \omega \geq 2$. Without loss of generality, we may assume that $\omega \leq \alpha$.

Let $K \in \mathcal{K}(\text{Forb}(H))$ be a p -core CRG and denote $g = g_K(p)$. By Lemma 6, any edge between vertices of different colors must be gray. Since H is a split graph, H would embed into any K with such a pair of vertices. So, the vertices in K are monochromatic. Let $K(\omega - 1, 0)$ denote the CRG with $\omega - 1$ white vertices and all edges gray. Let $K(0, \alpha - 1)$ denote the CRG with $\alpha - 1$ black vertices and all edges gray. So,

$$ed_{\text{Forb}(H)}(p) \leq \min \{g_{K(\omega-1,0)}(p), g_{K(0,\alpha-1)}(p)\} = \min \left\{ \frac{p}{\omega-1}, \frac{1-p}{\alpha-1} \right\}.$$

By virtue of the fact that a clique and independent set can intersect in at most one vertex, $h \leq \alpha + \omega \leq h + 1$.

¹We get the notation from Hung, Syslo, Weaver and West [9]. Barrett, Jepsen, Lang, McHenry, Nelson and Owens [8] define a clique-star, but it is a different type of graph.

Case 1. $\alpha + \omega = h + 1$.

In the case of $p = 1/2$, all p -core CRGs have all gray edges. Hence, we need only consider $K(\omega - 1, 0)$ and $K(0, \alpha - 1)$ and $ed_{\text{Forb}(H)}(1/2) = \min \left\{ \frac{1/2}{\omega-1}, \frac{1/2}{\alpha-1} \right\}$. Let $p \in (0, 1/2)$ and let v be a largest-weight vertex such that $x = \mathbf{x}(v)$. By Lemma 6(ii), every vertex is black and all edges are either white or gray. If v has $h - \omega$ neighbors, then $H \mapsto K$.

Thus, because x is the largest weight, Lemma 7(i) gives that

$$\begin{aligned} d_G(v) &\leq (h - \omega - 1)x \\ \frac{p - g}{p} + \frac{1 - 2p}{p}x &\leq (\alpha - 2)x \\ p - g &\leq (p\alpha - 1)x. \end{aligned}$$

If $p < 1/\alpha$, then $g > p \geq p/(\omega - 1)$. If $p \geq 1/\alpha$, then Corollary 8(i) gives that

$$\begin{aligned} p - g &\leq (p\alpha - 1)\frac{g}{1 - p} \\ p(1 - p) &\leq gp(\alpha - 1) \\ \frac{1 - p}{\alpha - 1} &\leq g, \end{aligned}$$

with equality if and only if K consists of $\alpha - 1$ black vertices. Hence equality requires that all edges of K be gray.

A similar argument, using Lemma 6(iii), shows that, for $p \in (1/2, 1)$, either $g > 1 - p \geq (1 - p)/(\alpha - 1)$ or $g \geq p/(\omega - 1)$, with equality if and only if K consists of $\omega - 1$ white vertices and all gray edges.

Case 2. $\alpha + \omega = h$.

Let $V(H) = A \cup W$ in which A is an independent set of size α and W is a clique of size ω . Similar to Case 1, $ed_{\text{Forb}(H)}(1/2) = \min \left\{ \frac{1/2}{\omega-1}, \frac{1/2}{\alpha-1} \right\}$. Next let $p \in (\frac{1}{2}, 1)$; hence all vertices are white and all edges are either black or gray.

Let v_1, \dots, v_ℓ be a maximal gray clique. That is, any edge between these vertices is gray and every vertex not in $\{v_1, \dots, v_\ell\}$ has at least one black neighbor in $\{v_1, \dots, v_\ell\}$. Let $x_i = \mathbf{x}(v_i)$ for $i = 1, \dots, \ell$ and let $X = \sum_{i=1}^{\ell} x_i$.

Each vertex in A is nonadjacent to some member of W , otherwise $\alpha + \omega = h + 1$. Consequently, $\ell \leq \omega - 1$ because H can be partitioned into ω

independent sets. Using Lemma 7(ii),

$$\begin{aligned} \sum_{i=1}^{\ell} [d_G(v_i) - X + x_i] &\leq (\ell - 1)(1 - X) \\ \sum_{i=1}^{\ell} \left[\frac{1-p-g}{1-p} + \frac{2p-1}{1-p} x_i - X + x_i \right] &\leq (\ell - 1)(1 - X) \\ \ell - \ell \frac{g}{1-p} + \frac{p}{1-p} X - \ell X &\leq (\ell - 1)(1 - X) \\ 1 - p - \ell g &\leq (1 - 2p)X. \end{aligned}$$

Hence, $g > \frac{1-p}{\ell} \geq \frac{1-p}{\omega-1} \geq \frac{1-p}{\alpha-1}$. From here on, we may assume $p \in (0, 1/2)$ and so all vertices are black and all edges are either white or gray.

Let $p \in \left(0, \frac{\omega-1}{h-1}\right]$. Let v be a vertex of largest weight $x = \mathbf{x}(v)$. Lemma 7(i) gives that

$$\begin{aligned} d_G(v) &\leq (h - \omega - 1)x \\ \frac{p-g}{p} + \frac{1-2p}{p}x &\leq (\alpha - 1)x \\ p - g &\leq (p(\alpha + 1) - 1)x. \end{aligned}$$

If $p < 1/(\alpha+1)$, then $g > p \geq p/(\omega-1)$. If $p \geq 1/(\alpha+1)$, then Corollary 8(i) gives that

$$p - g \leq (p(\alpha + 1) - 1) \frac{g}{1-p}.$$

Then,

$$g \geq \frac{1-p}{\alpha} \geq \frac{1 - \frac{\omega-1}{h-1}}{\alpha} = \frac{1}{h-1} = \frac{\frac{\omega-1}{h-1}}{\omega-1} \geq \frac{p}{\omega-1}.$$

Equality holds only if K has α black vertices and all edges gray. Since $V(H) = \bigcup_{a \in A} N[a]$, $H \mapsto K$ in that case.

Finally, we may assume that $p \in \left(\frac{\omega-1}{h-1}, \frac{1}{2}\right)$. We have to split into two cases according to the structure of H .

Case 2a. $\alpha + \omega = h$ and there exists an $c \leq \omega - 1$ such that H can be partitioned into c cliques and an independent set of $\alpha - c$ vertices.

Let v_1, \dots, v_ℓ be a maximal gray clique. That is, any edge between these vertices is gray and every vertex not in $\{v_1, \dots, v_\ell\}$ has at least one white neighbor in $\{v_1, \dots, v_\ell\}$. Let $x_i = \mathbf{x}(v_i)$ for $i = 1, \dots, \ell$ and let $X = \sum_{i=1}^{\ell} x_i$.

Using Lemma 7(i),

$$\begin{aligned} \sum_{i=1}^c [d_G(v_i) - X + x_i] &\leq (c-1)(1-X) \\ c \frac{p-g}{p} + \frac{1-p}{p} X - cX &\leq (c-1)(1-X) \\ p - cg &\leq (2p-1)X. \end{aligned}$$

Hence, $g > \frac{p}{c} \geq \frac{p}{\ell} \geq \frac{p}{\omega-1}$.

Which graphs are in Case 2, but not Case 2a? Since $\alpha + \omega = h$, every $w \in W$ has at least one neighbor in A . If any $a \in A$ has more than one neighbor in W , then we can greedily find at most $\omega - 1$ vertices in A such that the union of their neighborhoods is W . Such a graph would be in Case 2a.

So, the graphs, H with $\omega \leq \alpha$ that are in neither Case 1 nor Case 2a have the property that $N(w) \cap N(w') \cap A = \emptyset$ for all distinct $w, w' \in W$. This is exactly the case of a clique-star.

Case 2b. $\alpha + \omega = h$ and G is a clique-star.

Let $W = \{w_1, \dots, w_\omega\}$ such that w_i has $a_i + 1$ neighbors in A for $i = 1, \dots, \omega$ and there are a_0 isolated vertices.

Fact 11. *If $\omega \geq 2$ and H is a $(\omega; a_0, \dots, a_\omega)$ -clique-star and K is a black-vertex CRG such that either*

- *there exists a vertex with at least α gray neighbors, or*
- *there exist vertices v_1, \dots, v_ω such that*
 - *$\{v_1, \dots, v_\omega\}$ is a gray clique,*
 - *for $i = 1, \dots, \omega - 1$, v_i has $\alpha - 1$ gray neighbors, and*
 - *v_ω has at least $\lfloor (\alpha - \omega)/\omega \rfloor + \omega - 1$ gray neighbors (including $v_1, \dots, v_{\omega-1}$).*

Then, $H \mapsto K$.

Proof of Fact 11. If K has a vertex, v , with α gray neighbors, then W can be mapped to v whereas each member of $A = V(H) - W$ can be mapped to a different gray neighbor of v . Thus $H \mapsto K$. So, we may assume the maximum gray degree of K is at most $\alpha - 1$.

Our mapping is done recursively: Map w_ω and one of its neighbors to v_ω . Map its remaining A -neighbors ($a_\omega \leq \lfloor (\alpha - \omega)/\omega \rfloor$ of them) to each of a_ω gray neighbors of v_ω that are not in $\{v_1, \dots, v_{\omega-1}\}$.

Having embedded w_ω, \dots, w_{i+1} and each of their respective A -neighbors into a total of at most $\sum_{j=i+1}^{\omega} a_j$ vertices of K , we map w_i and one of its A -neighbors into v_i and its remaining a_i A -neighbors into arbitrary unused gray neighbors of v_i . After w_1 and its neighbors are mapped, we map the remaining a_0 isolated vertices arbitrarily into unused vertices of K .

This mapping can be accomplished because the fact that each of the v_i have at least $\alpha - 1$ gray neighbors ensures that, even at the last step, when w_1 and a neighbor is embedded, there are at least $\alpha - 1$ gray neighbors of v_1 . The number of gray neighbors of v_1 that were used are the $\omega - 1$ vertices v_i and at most $\sum_{j=2}^{\omega} a_j = \alpha - \omega - a_1 - a_0$ others, for a total of $\alpha - 1 - a_1 - a_0$. So, there are enough gray neighbors of v_1 to embed the a_1 neighbors of w_1 as well as the a_0 isolated vertices. Thus, $H \mapsto K$. \square

Fact 12. *Let $p \in (0, 1/2)$ and let K be a black-vertex CRG. If $g_K(p) \leq \min\{p/(\omega - 1), (1 - p)/(\alpha - 1)\}$, then either*

- *there exists a vertex with at least α gray neighbors, or*
- *there exist vertices v_1, \dots, v_{ω} such that*
 - *$\{v_1, \dots, v_{\omega}\}$ is a gray clique,*
 - *for $i = 1, \dots, \omega - 1$, v_i has $\alpha - 1$ gray neighbors, and*
 - *v_{ω} has at least $\lfloor (\alpha - \omega)/\omega \rfloor + \omega - 1$ gray neighbors (including $v_1, \dots, v_{\omega-1}$).*

Equality occurs if and only if $K \approx K(0, \alpha - 1)$.

Proof of Fact 12. Assume that no vertex has α neighbors. We find v_1, \dots, v_{ω} greedily. Choose v_1 to be a vertex of largest weight. Stop if $i = \omega$ or if $N_G(v_1) \cap \dots \cap N_G(v_i)$ is empty. Otherwise, let v_{i+1} be a vertex of largest weight in that set. We will show later that this process creates at least ω vertices.

First, we find the number of gray neighbors of v_1 , using the fact that x_1 is the largest weight.

$$|N_G(v_1)| \geq \left\lceil \frac{d_G(v_1)}{x_1} \right\rceil \geq \frac{p - g}{px_1} + \frac{1 - 2p}{p}.$$

Using Corollary 8(i), we have that $x_1 \leq g/(1 - p)$ and so

$$|N_G(v_1)| \geq \frac{1 - p - g}{g} \geq \alpha - 2.$$

Equality only occurs if $g = (1 - p)/(\alpha - 1)$ there are $\alpha - 1$ vertices, all of weight $1/(\alpha - 1)$, thus $K \approx K(0, \alpha - 1)$. So, we may assume $|N_G(v_1)| \geq \alpha - 1$.

For $i \in \{2, \dots, \omega - 1\}$, we let $X = \sum_{j=1}^i x_j$ and consider the gray neighborhood of v_i , excluding $\{v_1, \dots, v_{i-1}\}$. Its total weight is:

$$(5) \quad d_G(v_i) - (X - x_i) - \sum_{j=1}^{i-1} \mathbf{x}(N_W(v_j)) = \frac{p - ig}{p} + \frac{1 - 2p}{p}X > 0,$$

because $i \leq \omega - 1$, $g \leq p/(\omega - 1)$, $p < 1/2$ and $X > x_i > 0$. Thus, v_{i+1} can be obtained.

We use these calculations to obtain the size of $N_G(v_i)$ for $i = 2, \dots, \omega - 1$. First note that v_i has $i - 1$ gray neighbors among $\{v_1, \dots, v_{i-1}\}$ and that every vertex that is a gray neighbor of each of v_1, \dots, v_i has weight at most x_i . As to the remaining vertices, partition $N_G(v_i)$ according to the least index j for which the vertex is adjacent to v_j via a white edge. By the choice

of v_1, \dots, v_i , such a vertex has weight at most $x_j = \mathbf{x}(v_j)$. Consequently, we have a lower bound for $|N_G(v_i)|$:

$$|N_G(v_i)| \geq (i-1) + \left\lceil \frac{\mathbf{x}(N_G(v_1) \cap \dots \cap N_G(v_i))}{x_i} \right\rceil + \sum_{j=1}^{i-1} \left\lceil \frac{\mathbf{x}(N_G(v_i) \cap N_W(v_j) \cap \{N_G(v_1) \cap \dots \cap N_G(v_{j-1})\})}{x_j} \right\rceil.$$

We can drop the ceilings to obtain the lower bound

$$|N_G(v_i)| \geq (i-1) + \frac{1}{x_i} \mathbf{x}(N_G(v_1) \cap \dots \cap N_G(v_i)) + \sum_{j=1}^{i-1} \frac{1}{x_j} \mathbf{x}(N_G(v_i) \cap N_W(v_j) \cap \{N_G(v_1) \cap \dots \cap N_G(v_{j-1})\}).$$

Now we look at the coefficients $\frac{1}{x_1} < \frac{1}{x_2} < \dots < \frac{1}{x_i}$. The total weight of gray neighbors with coefficient $\frac{1}{x_1}$ is at most $\mathbf{x}(N_W(v_1))$. The total weight of gray neighbors with coefficient $\frac{1}{x_1}$ or $\frac{1}{x_2}$ is at most $\mathbf{x}(N_W(v_1)) + \mathbf{x}(N_W(v_2))$ and so on.

$$|N_G(v_i)| \geq (i-1) + \sum_{j=1}^{i-1} \frac{1}{x_j} \mathbf{x}(N_W(v_j)) + \frac{\mathbf{x}(N_G(v_i)) - (X - x_i) - \sum_{j=1}^{i-1} \mathbf{x}(N_W(v_j))}{x_i}$$

and observe that inequality (5) shows that the last numerator is nonnegative.

Using similar computations as before,

$$|N_G(v_i)| \geq (i-1) + \sum_{j=1}^{i-1} \frac{1}{x_j} \left(\frac{g}{p} - \frac{1-p}{p} x_j \right) + \frac{1}{x_i} \left(\frac{p-g}{p} + \frac{1-2p}{p} x_i - (X - x_i) - \sum_{j=1}^{i-1} \left(\frac{g}{p} - \frac{1-p}{p} x_j \right) \right).$$

After some simplification

$$(6) \quad |N_G(v_i)| \geq \frac{g}{p} \sum_{j=1}^{i-1} \frac{1}{x_j} - \frac{1-2p}{p} (i-1) + \frac{p-ig}{px_i} + \frac{1-2p}{p} \left(\frac{X}{x_i} \right).$$

Using Janson's inequality and the fact that $X - x_i \geq \frac{i-1}{i} X$, we see that

$$\sum_{j=1}^{i-1} \frac{1}{x_j} \geq \frac{i-1}{(X - x_i)/(i-1)} \geq \frac{i(i-1)}{X}.$$

So, we return to (6) and then the fact that $x_i \leq X/i$:

$$\begin{aligned} |N_G(v_i)| &\geq \frac{g}{p} \left(\frac{i(i-1)}{X} \right) - \frac{1-2p}{p}(i-1) + \frac{p-ig+(1-2p)X}{px_i} \\ &\geq \frac{gi(i-1)}{pX} - \frac{1-2p}{p}(i-1) + \frac{p-ig+(1-2p)X}{p(X/i)} \\ &= \frac{i(p-g)}{pX} + \frac{1-2p}{p}. \end{aligned}$$

Using the fact that $X \leq ig/(1-p)$, we see that

$$|N_G(v_i)| \geq \frac{1-p-g}{g} \geq \alpha - 2.$$

Equality only occurs if $g = (1-p)/(\alpha-1)$ and $K \approx K(0, \alpha-1)$.

Finally, we try to determine the number of vertices adjacent to v_ω via a gray edge. We only need $|N_G(v_\omega)| \geq \lfloor \alpha/\omega \rfloor + \omega - 2$ in order to finish the proof. First, note that the very existence of v_ω ensures that $|N_G(v_\omega)| \geq \omega - 1$. Thus, we may assume that $\alpha \geq 2\omega$.

Second, suppose that $\omega \geq 3$. Recalling that $d_G(v_\omega) = \frac{p-g}{p} + \frac{1-2p}{p}x_\omega$ and $x_1 \leq \frac{g}{1-p}$, the pigeonhole principle gives that for $v = v_\omega$ (indeed, for any vertex v),

$$\begin{aligned} |N_G(v)| &\geq \left\lceil \frac{p-g}{p} \cdot \frac{1-p}{g} \right\rceil \\ &\geq \begin{cases} \left\lceil \frac{p-\frac{p}{\omega-1}}{p} \cdot \frac{1-p}{p/(\omega-1)} \right\rceil, & \text{if } p \leq \frac{\omega-1}{h-2}; \\ \left\lceil \frac{p-\frac{1-p}{\alpha-1}}{p} \cdot \frac{1-p}{(1-p)/(\alpha-1)} \right\rceil, & \text{if } p \geq \frac{\omega-1}{h-2}. \end{cases} \\ &\geq \left\lceil (\alpha-1) \frac{\omega-2}{\omega-1} \right\rceil. \end{aligned}$$

Let $\alpha = q\omega + r$ with $0 \leq r \leq \omega - 1$ and note that since $\alpha \geq 2\omega$, $q \geq 2$. Hence,

$$\begin{aligned} |N_G(v)| &\geq \left\lceil (\alpha-1) \frac{\omega-2}{\omega-1} \right\rceil \\ &= \left\lceil q(\omega-2) + \frac{(q+r-1)(\omega-2)}{\omega-1} \right\rceil \\ &\geq q(\omega-2) + 1 \\ &= q + \omega - 2 + (q-1)(\omega-3) \end{aligned}$$

Since $q = \lfloor \alpha/\omega \rfloor$, we may conclude that $|N_G(v)| \geq \lfloor \alpha/\omega \rfloor + \omega - 2$, just as desired.

Third, let $\omega = 2$; i.e., H is a double-star (possibly with isolated vertices). Recall that $\alpha \geq 2\omega = 4$. Our goal is to show that $|N_G(v_2)| \geq \lfloor \alpha/\omega \rfloor + \omega - 2 = \lfloor \alpha/2 \rfloor$. The computations are, by now, routine. We use the fact that

$x_2 \geq d_G(v_1)/(\alpha - 1)$ and $x_1 \leq g/(1 - p)$.

$$\begin{aligned}
|N_G(v_2)| &\geq \left\lceil \frac{d_G(v_2)}{x_1} \right\rceil \geq \left\lceil \frac{1}{x_1} \left(\frac{p-g}{p} + \frac{1-2p}{p} x_2 \right) \right\rceil \\
&\geq \left\lceil \frac{1}{x_1} \left(\frac{p-g}{p} + \frac{1-2p}{p} \cdot \frac{d_G(v_1)}{\alpha-1} \right) \right\rceil \\
&\geq \left\lceil \frac{p-g}{px_1} \left(1 + \frac{1-2p}{p(\alpha-1)} \right) + \left(\frac{1-2p}{p} \right)^2 \frac{1}{\alpha-1} \right\rceil \\
&\geq \left\lceil \frac{(p-g)(1-p)}{pg} \left(\frac{p(\alpha-3)+1}{p(\alpha-1)} \right) + \left(\frac{1-2p}{p} \right)^2 \frac{1}{\alpha-1} \right\rceil.
\end{aligned}$$

Recalling that, in the case of $\omega = 2$, $g \leq \min \{p, (1-p)/(\alpha-1)\}$,

$$\begin{aligned}
|N_G(v_2)| &\geq \begin{cases} \left\lceil \left(\frac{1-2p}{p} \right)^2 \frac{1}{\alpha-1} \right\rceil, & \text{if } p \leq 1/\alpha; \\ \left\lceil \frac{p\alpha-1}{p} \left(\frac{p(\alpha-3)+1}{p(\alpha-1)} \right) + \left(\frac{1-2p}{p} \right)^2 \frac{1}{\alpha-1} \right\rceil, & \text{if } p \geq 1/\alpha. \end{cases} \\
&= \begin{cases} \left\lceil \left(\frac{1-2p}{p} \right)^2 \frac{1}{\alpha-1} \right\rceil, & \text{if } p \leq 1/\alpha; \\ \left\lceil \alpha - 2 - \frac{(1-2p)}{p(\alpha-1)} \right\rceil, & \text{if } p \geq 1/\alpha. \end{cases}
\end{aligned}$$

In each case, the smallest value of the expression occurs when $p = 1/\alpha$, giving

$$|N_G(v_2)| \geq \left\lceil \alpha - 3 + \frac{1}{\alpha-1} \right\rceil = \alpha - 2 = \left\lfloor \frac{\alpha}{2} \right\rfloor + \left(\left\lceil \frac{\alpha}{2} \right\rceil - 2 \right).$$

This is at least $\lfloor \alpha/2 \rfloor$ since $\alpha \geq 4$. This concludes the proof of Fact 12. \square

Summarizing, if $H \not\rightarrow K$, then either $g \geq p/(\omega - 1)$ or $g \geq (1-p)/(\alpha - 1)$, with equality if and only if either $K \approx K(0, h - \omega)$ or $K \approx K(h - \alpha, 0)$. This concludes the proof of Theorem 1.

5.2. Examples of split graphs. Items (i) and (ii) in Corollary 13 were proven in [7].

Corollary 13. *Let H be a graph on h vertices.*

- (i) *If $H \approx K_a + E_b$, then $ed_{\text{Forb}(H)}(p) = \min \left\{ \frac{p}{a-1}, \frac{1-p}{b} \right\}$.*
- (ii) *If H is a star (i.e., $H \approx E_{h-1} \vee K_1$), then $ed_{\text{Forb}(H)}(p) = \min \left\{ p, \frac{1-p}{h-2} \right\}$.*
- (iii) *If H is a double-star (i.e., there are adjacent vertices u and v to which every other vertex is adjacent to exactly one), then $ed_{\text{Forb}(H)}(p) = \min \left\{ p, \frac{1-p}{h-3} \right\}$.*

6. $\text{Forb}(H_9)$

Marchant and Thomason [11] give the example of $\mathcal{H} = \text{Forb}(C_6^*)$, where C_6^* is a 6-cycle with an additional diagonal edge, such that $ed_{\mathcal{H}}(p)$ is not determined by CRGs with all gray edges. More precisely, they prove that

$$ed_{\text{Forb}(C_6^*)}(p) = \min \left\{ \frac{p}{1+2p}, \frac{1-p}{2} \right\}.$$

The CRG which corresponds to $g_K(p) = (1-p)/2$ is $K(0, 2)$. The CRG, K , which has $g_K(p) = p/(1+2p)$ for $p \in [0, 1/2]$ consists of three vertices: two black vertices connected via a white edge and a white vertex. The remaining two edges are gray.

The graph H_9 , shown in Figure 1.2 and cited in [7], generates a hereditary property $\mathcal{H} = \text{Forb}(H_9)$ such that $d_{\mathcal{H}}^*$ cannot be determined by CRGs of the form $K(a, c)$. Note that $d_{\text{Forb}(C_6^*)}$ can be determined by such CRGs, but the part of the function for $p \in (0, 1/2)$ cannot.

6.1. Proof of Theorem 3. Upper bound. We know that $\chi(H_9) = 4$ so let $K^{(1)} = K(3, 0)$ where $g_{K^{(1)}}(p) = p/3$. We also know that $\chi(\overline{H_9}) = 3$ so let $K^{(4)} = K(0, 2)$ where $g_{K^{(4)}}(p) = (1-p)/2$. In [7], another CRG in $\mathcal{K}(\text{Forb}(H_9))$ is given, call it $K^{(2)}$. It consists of 4 white vertices, one black edge and 5 gray edges. It has edit distance function $g_{K^{(2)}}(p) = \min\{p/3, p/(2+2p)\}$.

There is a CRG with a smaller g function. We call it $K^{(3)}$, it consists of 5 white vertices, two disjoint black edges and the remaining 8 edges gray. The function $g_{K^{(2)}}(p)$ can be computed by use of Theorem 10. In the setup of that theorem, $K^{(3)}$ has 3 components. Since the components have g functions either p (for the solitary white vertex) or $\min\{p, 1/2\}$ (for each of the other two components), the theorem gives that

$$g_{K^{(3)}}(p)^{-1} = p^{-1} + 2(\min\{p, 1/2\})^{-1} = \min\{p/3, p/(1+4p)\}.$$

It is easy to see that $H_9 \not\mapsto K^{(1)}$ and $H_9 \not\mapsto K^{(4)}$. In [7], it was shown that $H_9 \not\mapsto K^{(2)}$. To finish the upper bound, it remains to show that $H_9 \not\mapsto K^{(3)}$. Let v_0 be the isolated vertex, $\{v_1, w_1\}$ be a black edge and $\{v_2, w_2\}$ be a black edge.

First, we show that no component of $K^{(3)}$ can have 4 vertices from H_9 . Since there are no independent sets of size 4 and no induced stars of size 4, the only way to have a component of size 4 is to have an induced copy of C_4 in the component consisting of, say, $\{v_2, w_2\}$. It is not difficult to see that deleting two vertices from the set $\{0, 3, 6\}$ yields a C_4 -free graph. So, any C_4 contains exactly two members of $\{0, 3, 6\}$. Without loss of generality, the induced C_4 is $\{1, 3, 6, 8\}$. But the graph induced by $\{0, 2, 4, 5, 7\}$ induces a C_5 , which cannot be mapped into the sub-CRG induced by $\{v_0, v_1, w_1\}$. Therefore, if H_9 were to map to $K^{(3)}$, each component must contain exactly 3 vertices. First we map to v_0 . The only independent sets of size 3 are $\{1, 4, 7\}$ and $\{2, 5, 8\}$. Without loss of generality, assume the former. Second,

we consider the graph induced by $\{0, 2, 3, 5, 6, 8\}$. Any partition of these vertices into two subsets of 3 vertices either has a triangle or a copy of \overline{P}_3 , neither of which maps into $\{v_1, w_1\}$ or $\{v_2, w_2\}$. So, these six vertices cannot be mapped into $\{v_1, w_1, v_2, w_2\}$. Hence $H_9 \not\rightarrow K^{(3)}$.

The CRGs $K^{(1)}$, $K^{(3)}$ and $K^{(4)}$ give an upper bound of $\min\left\{\frac{p}{3}, \frac{p}{1+4p}, \frac{1-p}{2}\right\}$.

Lower bound, for $p \leq 1/2$. Let K be a p -core such that $H_9 \not\rightarrow K$. If K has at least 2 white vertices, then it has no black vertices because $H_9 \mapsto K(2, 1)$. (The independent sets are $\{1, 4, 7\}$ and $\{2, 5, 8\}$ and the clique is $\{0, 3, 6\}$.) So, in this case $g_K(p) \geq p/3$ with equality if and only if $K \approx K(3, 0)$.

If K has exactly one white vertex, then there is no gray edge among the black vertices because $H_9 \mapsto K(1, 2)$. (The independent set is $\{2, 7\}$ and the cliques are $\{0, 1, 8\}$ and $\{3, 4, 5, 6\}$.) Let w be the white vertex and $K' = K - \{w\}$ and $k' = |V(K')|$. Since K' is a clique with all black vertices and all white edges, Proposition 8 from [12] gives that, for $p \in (0, 1/2)$, $g_{K'}(p) = p + \frac{1-2p}{k'} > p$. By Theorem 10, $g_K(p) > 1/(1/p + 1/p) = p/2$, which is strictly larger than $ed_{\text{Forb}(H_9)}(p)$ for $p \in (0, 1/2]$.

If K has no white vertices, then let v_0 be the vertex with largest weight and let v_1 be a vertex in the gray neighborhood of v_0 . Let $x_0 = \mathbf{x}(v_0)$ and $x_1 = \mathbf{x}(v_1)$. Since K can have no gray triangles (H_9 can be partitioned into 3 cliques), $d_G(v_0) + d_G(v_1) \leq 1$.

$$\begin{aligned} 1 &\geq d_G(v_0) + d_G(v_1) \\ &\geq 2\frac{p-g}{p} + \frac{1-2p}{p}(x_0 + x_1) \\ g &\geq \frac{p}{2} + \frac{1-2p}{2}(x_0 + x_1) > \frac{p}{2}. \end{aligned}$$

Summarizing, if $p \leq 1/2$ and K is a p -core such that $H \not\rightarrow K$, then $g_K(p) \geq p/3$ with equality only if $K \approx K(3, 0)$.

Lower bound, for $p \geq 1/2$. Let K be a p -core such that $H_9 \not\rightarrow K$. If K has at least 2 black vertices, then there are no white vertices because $H_9 \mapsto K(1, 2)$ and so $g_K(p) \geq (1-p)/2$ with equality if and only if $K \approx K(0, 2)$.

If K has exactly one black vertex, then there is no gray edge among the white vertices because $H_9 \mapsto K(2, 1)$. Let b be the black vertex and $K' = K - \{b\}$ and $k' = |V(K')|$. Similar to the above, Proposition 8 from [12] can be used to show that, for $p \in (1/2, 1)$, $g_{K'}(p) = 1 - p + \frac{2p-1}{k'} > 1 - p$. By Theorem 10, $g_K(p) > (1-p)/2$, which is strictly larger than $ed_{\text{Forb}(H_9)}(p)$ for $p \in [1/2, 1)$.

From now on, we will assume that K has only white vertices and, since it is p -core for $p \geq 1/2$, all edges are black or gray. Fact 14 and Fact 15 establish some of the structural theorems.

Fact 14. *Let $p \in [1/2, 1)$ and K be a p -core CRG with white vertices and black or gray edges. Let v and v' be vertices connected by a gray edge. Then, $N_G(v) \cap N_G(v')$ has at most two vertices.*

Proof. If $N_G(v) \cap N_G(v')$ has three vertices, then map H_9 vertices 0, 3 and 6 to each of them, map $\{1, 4, 7\}$ to v and $\{2, 5, 8\}$ to v' . This is a map demonstrating that $H_9 \mapsto K$. \square

Fact 15. *Let $p \in [1/2, 1)$ and K be a p -core CRG with white vertices and black or gray edges. Let v_0 be a vertex of largest weight and v_1 be a vertex that has largest weight among those in $N_G(v_0)$. Then, either $N_G(v_0) \cap N_G(v_1)$ has exactly two vertices or $g_K(p) > (1-p)/2$ or $g_K(p) \geq p/3$ with equality if and only if $K \approx K(3, 0)$.*

Proof. Let $g = g_K(p)$. If the statement of Fact 15 is not true, then $N_G(v_0) \cap N_G(v_1)$ has at most one vertex which, by the choice of v_1 , has weight at most $\mathbf{x}(v_1)$ and, by inclusion-exclusion, has weight at least $d_G(v_0) + d_G(v_1) - 1$. Therefore,

$$\begin{aligned} \mathbf{x}(v_1) &\geq d_G(v_0) + d_G(v_1) - 1 \\ &\geq 2\frac{1-p-g}{1-p} + \frac{2p-1}{1-p}(\mathbf{x}(v_0) + \mathbf{x}(v_1)) - 1 \\ (7) \quad g &\geq \frac{1-p}{2} + \frac{2p-1}{2}\mathbf{x}(v_0) - \frac{2-3p}{2}\mathbf{x}(v_1). \end{aligned}$$

If $p \geq 2/3$, then $g > (1-p)/2$. If $p < 2/3$, then use $\mathbf{x}(v_1) \leq \mathbf{x}(v_0)$ in (7).

$$(8) \quad g \geq \frac{1-p}{2} + \frac{5p-3}{2}\mathbf{x}(v_1)$$

If $p \geq 3/5$, then $g > (1-p)/2$. If $p < 3/5$, then use the fact that Corollary 8(ii) gives $\mathbf{x}(v_0) \leq g/p$, which we use in (8).

$$\begin{aligned} g &\geq \frac{1-p}{2} + \frac{5p-3}{2}\mathbf{x}(v_1) \geq \frac{1-p}{2} + \frac{5p-3}{2}\left(\frac{g}{p}\right) \\ g &\geq \frac{p}{3}. \end{aligned}$$

It is easy to see that equality can only occur if $\mathbf{x}(v_2) = \mathbf{x}(v_1) = g/p = 1/3$ and their common gray neighborhood is a vertex of weight $1 - 2g/p = 1/3$. \square

Given Fact 14 and Fact 15, we can identify v_0 , a vertex of maximum weight, v_1 a vertex of maximum weight among those in $N_G(v_0)$ and $\{v_2, w_2\} = N_G(v_0) \cap N_G(v_1)$. Without loss of generality, let $\mathbf{x}(v_2) \geq \mathbf{x}(w_2)$. For ease of notation, let $x_i = \mathbf{x}(v_i)$ for $i = 0, 1, 2$. If $N_G(v_0) \cap N_G(v_2) - \{v_1\}$ is nonempty, then let its unique vertex be denoted w_1 . (Uniqueness is a consequence of Fact 14.)

Case 1. The vertex w_1 does not exist.

Most of our observations come from inclusion-exclusion: $|A| + |B| = |A \cup B| + |A \cap B|$. Inequality (9) comes from the fact that $N_G(v_0) \cap N_G(v_1) = \{v_2, w_2\}$. Inequality (10) comes from the fact that $N_G(v_0) \cap N_G(v_2) = \{v_1\}$. Observe that $\mathbf{x}(w_2) \leq x_2$, hence,

$$(9) \quad d_G(v_0) + d_G(v_1) \leq 1 + 2x_2$$

$$(10) \quad d_G(v_0) + d_G(v_2) \leq 1 + x_1.$$

Solve for x_2 in each case, recalling that Lemma 7(ii) gives that $d_G(v_2) = \frac{1-p-g}{1-p} + \frac{2p-1}{1-p}x_2$. Inequality (9) gives a lower bound for x_2 and inequality (10) gives an upper bound:

$$\frac{1}{2}(d_G(v_0) + d_G(v_1) - 1) \leq x_2 \leq \frac{1-p}{2p-1} \left(1 + x_1 - d_G(v_0) - \frac{1-p-g}{1-p} \right).$$

Some simplification gives

$$\begin{aligned} 2g &\geq d_G(v_0) + (2p-1)d_G(v_1) - 2(1-p)x_1 - 2p + 1 \\ &\geq 2p \frac{1-p-g}{1-p} + \frac{2p-1}{1-p}x_0 + \frac{2p^2-1}{1-p}x_1 - 2p + 1 \\ g &\geq \frac{1-p}{2} + \frac{2p-1}{2}x_0 + \frac{2p^2-1}{2}x_1. \end{aligned}$$

If $2p^2 - 1 > 0$ (i.e, $p > 1/\sqrt{2}$), then $g > (1-p)/2$. Otherwise, we use the bound $x_1 \leq x_0$.

$$\begin{aligned} g &\geq \frac{1-p}{2} + \frac{2p-1}{2}x_0 + \frac{2p^2-1}{2}x_0 \\ &\geq \frac{1-p}{2} + (p^2 + p - 1)x_0. \end{aligned}$$

If $p^2 + p - 1 > 0$ (i.e, $p > (\sqrt{5} - 1)/2$), then $g > (1-p)/2$. Otherwise, we use the bound from Corollary 8(ii) that $x_0 \leq g/p$.

$$\begin{aligned} g &\geq \frac{1-p}{2} + (p^2 + p - 1)x_0 \\ &\geq \frac{1-p}{2} + (p^2 + p - 1)\frac{g}{p} \\ &\geq \frac{p}{2(1+p)}. \end{aligned}$$

Equality occurs only if $x_0 = x_1 = g/p$ and $\mathbf{x}(w_2) = x_2 = 1/2 - g/p$. This is precisely the CRG denoted $K^{(2)}$.

Case 2. The vertex w_1 exists.

Inequality (11) comes from the fact that $N_G(v_0) \cap N_G(v_1) = \{v_2, w_2\}$ and $\mathbf{x}(w_2) \leq \mathbf{x}(v_2) = x_2$. Inequality (12) comes from the fact that $N_G(v_0) \cap$

$N_G(v_2) = \{v_1, w_1\}$ and $\mathbf{x}(w_1) \leq \mathbf{x}(v_1) = x_1$. Observe that $\mathbf{x}(w_2) \leq x_2$ and $\mathbf{x}(w_1) \leq x_1$, hence,

$$(11) \quad d_G(v_0) + d_G(v_1) \leq 1 + 2x_2$$

$$(12) \quad d_G(v_0) + d_G(v_2) \leq 1 + 2x_1.$$

Adding (11) and (12) gives

$$(13) \quad \begin{aligned} 2d_G(v_0) + d_G(v_1) + d_G(v_2) &\leq 2 + 2(x_1 + x_2) \\ 2d_G(v_0) - \frac{2g}{1-p} &\leq \frac{3-4p}{1-p}(x_1 + x_2). \end{aligned}$$

If $p \geq 3/4$, then (13) gives that $2d_G(v_0) - \frac{2g}{1-p} \leq 0$. Consequently, $g > (1-p)/2$. So, we assume $p < 3/4$.

Next, we use Fact 16 to conclude that v_0 is the only common gray neighbor of v_1 and v_2 .

Fact 16. *Let $p \geq 1/2$ and K be a p -core with white vertices and black or gray edges. Let $a_0, a_1, a_2, b_0, b_1, b_2 \in V(K)$ such that $\{a_0, a_1, a_2\}$ is a gray triangle and $\{b_i, a_j\}$ is a gray edge as long as i and j are distinct. Then, $H_9 \mapsto K$.*

Proof. The following map shows the embedding:

$$\begin{array}{lll} 2, 7 \mapsto a_0 & 1, 5 \mapsto a_1 & 4, 8 \mapsto a_2 \\ 0 \mapsto b_0 & 3 \mapsto b_1 & 6 \mapsto b_2. \end{array}$$

□

If v_1 and v_2 have a gray neighbor in K other than v_0 , call it w_0 and observe that by setting $a_i := v_i$ and $b_i := w_i$ for $i = 0, 1, 2$, Fact 16 would imply that $H_9 \mapsto K$.

Since v_0 is the only common gray neighbor of v_1 and v_2

$$(14) \quad \begin{aligned} d_G(v_1) + d_G(v_2) &\leq 1 + x_0 \\ \frac{2p-1}{1-p}(x_1 + x_2) &\leq 1 + x_0 - 2\frac{1-p-g}{1-p}. \end{aligned}$$

Inequality (13) gives a lower bound for $x_1 + x_2$ and inequality (14) gives an upper bound. Recall that Lemma 7(ii) gives that $d_G(v) = \frac{1-p-g}{1-p} + \frac{2p-1}{1-p}\mathbf{x}(v)$ for any vertex $v \in V(K)$. Recall that we assume $p \leq 3/4$.

$$\frac{1-p}{3-4p} \left(2d_G(v_0) - \frac{2g}{1-p} \right) \leq x_1 + x_2 \leq \frac{1-p}{2p-1} \left(1 + x_0 - 2\frac{1-p-g}{1-p} \right).$$

Some simplification gives

$$2(2p-1)((1-p)d_G(v_0) - g) \leq (3-4p)((1-p)(1+x_0) - 2(1-p-g))$$

and so

$$g \geq \frac{1-p}{2} + \frac{4p^2 - p - 1}{2}x_0.$$

If $4p^2 - p - 1 > 0$ (i.e, $p > (\sqrt{17} + 1)/8$), then $g > (1 - p)/2$. Otherwise, we use the bound $x_0 \leq g/p$ from Corollary 8(ii).

$$\begin{aligned} g &\geq \frac{1-p}{2} + \frac{4p^2 - p - 1}{2} \left(\frac{g}{p}\right) \\ &\geq \frac{p}{1+4p}. \end{aligned}$$

Equality occurs only if $x_0 = g/p$, $x_1 = x_2 = \frac{p}{1+4p}$ and $\mathbf{x}(w_i) = x_i$ for $i = 1, 2$. This is precisely the CRG denoted $K^{(3)}$.

Therefore, for $p \in [1/2, 1]$ and in each case, $g \geq \min \{p/(1+4p), (1-p)/2\}$. Combining this with the fact that for $p \in [0, 1]$ that $g \geq p/3$. This concludes the proof of the lower bound. Consequently, $ed_{\text{Forb}(H_9)}(p) = \min \{p/3, p/(1+4p), (1-p)/2\}$. This concludes the proof of Theorem 3.

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