# The Number of Arcs of Strongly Connected Oriented Graphs with Two Noncritical Vertices<sup>\*</sup>

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**Abstract** It is proved that a strongly connected oriented graph D with  $n \ge 4$  vertices and at least  $\binom{n-1}{2} + 3$  arcs has two distinct vertices  $u^*, v^*$  such that both  $D - u^*$  and  $D - v^*$  are strongly connected. The examples show that the above lower bound on the number of arcs is sharp.

Keywords digraph, strongly connected subdigraph, critical vertex Chinese Library Classification 0157.5

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## 含有两个非临界点的强连通定向图的弧数

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    摘要 证明顶点数为 n ≥ 4, 弧数为 m ≥ (<sup>n-1</sup>) + 3 的强连通定向图 D 中存在两点 u<sup>*</sup>、
    v<sup>*</sup>, 使得 D - u<sup>*</sup> 和 D - v<sup>*</sup> 都是强连通的,并用例子说明这里所给的关于弧数的下界是紧的.
    关键词 有向图,强连通子图,临界点
    中图分类号 O157.5
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#### **1** Terminology and introduction

For graph-theoretical terminology and notation not defined here we follow [1]. We only consider finite (di)graphs without loops and multiple edges (arcs). A digraph H is a subdigraph of a digraph D (written  $H \subseteq D$ ) if  $V(H) \subseteq V(D)$ ,  $A(H) \subseteq A(D)$  and every arc in A(H) has both end-vertices in V(H). When  $H \subseteq D$  but  $H \neq D$ , we call H a proper subdigraph of D. A spanning subdigraph of D is a subdigraph H with V(H) = V(D). Suppose that U is a nonempty subset of V = V(D). The subdigraph of D whose vertex set is U and whose arc set is the set of those arcs of D that have both end-vertices in U is called the subdigraph of D. The induced by U and is denoted by D[U]; we say that D[U] is an induced subdigraph of D. The induced subdigraph  $G[V \setminus U]$  is denoted by D - U; it is the

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subdigraph obtained from D by deleting the vertices in U together with their incident arcs. If  $U = \{v\}$  we write D - v for  $D - \{v\}$ .

A digraph D is strongly connected (or, just, strong) if every vertex of D is reachable from every other vertex of D. We define a digraph with one vertex to be strong. A vertex v in a strong digraph D is said to be noncritical if the digraph D - v is also strong. Similar notions can be introduced for an undirected graph G. It is well known that a connected undirected graph G with at least two vertices contains a spanning tree T, whence at least two noncritical vertices (in particular, any leaf in T is noncritical both for T and G). So we have the following statement.

**Observation 1.1** Every connected undirected graph on at least 2 vertices has at least two noncritical vertices.

An oriented graph is a digraph with no cycle of length two. An out-arborescence rooted at s is an oriented tree T such that  $s \in V(T)$  and the in-degree of every vertex  $x \in V(T) \setminus \{s\}$  is 1. A tournament is an oriented graph where every pair of distinct vertices are adjacent. We have a similar result to Observation 1.1 for tournaments.

**Theorem 1.1**<sup>[1]</sup> Every strong tournament on at least 4 vertices contains at least two noncritical vertices.

Consider the directed cycle  $C_n$  with  $n \ge 3$  vertices. It is easy to see that  $C_n$  is strong, but contains no noncritical vertex. So Observation 1.1 can not be extended to general digraphs. In 1999, Schwarz<sup>[2]</sup> conjectured that every strong digraph with  $n \ge 3$  vertices and  $m \ge {n \choose 2} + 1$  arcs contains at least one noncritical vertex. In the same year, London<sup>[3]</sup> proved this conjecture. Later, Aharoni and Berger<sup>[4]</sup> showed that the lower bound on the number of arcs can be sharpened from  ${n \choose 2} + 1$  to  ${n-1 \choose 2} + 4$ .

**Theorem 1.2**<sup>[4]</sup> If a strong digraph D with  $n \ge 4$  vertices has at least  $\binom{n-1}{2} + 4$  arcs, it has at least one noncritical vertex.

For oriented graphs, the above bound on the number of arcs can be improved to  $\binom{n-2}{2} + 5$ .

**Theorem 1.3**<sup>[5]</sup> If a strong oriented graph D with  $n \ge 4$  vertices has at least  $\binom{n-2}{2} + 5$  arcs, it has at least one noncritical vertex.

In this paper, we will show that  $\binom{n-1}{2} + 3$  arcs in a strong oriented graph with  $n \ge 4$  vertices can guarantee the existence of two noncritical vertices. Clearly, Theorem 1.1 is an immediate consequence of this result.

#### 2 Main result

We start with the notion of maximal strong proper subdigraphs and a characterization of such subdigraphs. A strong proper subdigraph H of a strong digraph D is maximal if any strong subdigraph in D containing H coincides either with H or with D.

**Lemma 2.1**<sup>[6]</sup> Let H be a strong proper subdigraph of a strong digraph D, and let H' = D - V(H). Then H is a maximal strong proper subdigraph in D iff the following three conditions are satisfied:

(1) there is a vertex  $w_{in}$  in H' such that any arc going from V(H) to V(H') enters  $w_{in}$ ;

(2) there is a vertex  $w_{out}$  in H' such that any arc going from V(H') to V(H) leaves from  $w_{out}$ ;

(3) the distance between  $w_{in}$  and  $w_{out}$  is one less than the order of H'.

The main result of this paper is the following.

**Theorem 2.1** If a strong oriented graph D with  $n \ge 4$  vertices has at least  $\binom{n-1}{2} + 3$  arcs, it has at least two noncritical vertices.

**Proof** By induction on n. The statement is clearly true for n = 4. Suppose, then, that  $n \ge 5$ . By Theorem 1.3, D has a vertex  $v^*$  such that  $D - v^*$  is strong. We will show that D has a vertex  $u^* \ne v^*$  such that  $D - u^*$  is also strong. Let H be a maximal strong proper subdigraph of D, which contains the vertex  $v^*$ , and let H' = D - V(H). Since the subdigraph induced by  $\{v^*\}$  is strong, such a subdigraph H exists. By Lemma 2.1, H' has a Hamilton path  $x_1x_2 \dots x_k$  such that any arc going from V(H) to V(H') enters  $x_1$ , any arc going from V(H') to V(H) leaves from  $x_k$ , and there is no arc of the form  $x_ix_j$ , where  $i \le j-2$ .

If k = 1, we are done since  $H = D - x_1$  is strong and  $x_1 \neq v^*$ .

If  $k \ge 3$ , we construct the digraph D' from D by removing the vertex  $x_1$  and adding all the arcs  $wx_2$  where w is an in-neighbor of  $x_1$  in H, that is,

$$D' = (D - x_1) + \{ wx_2 : wx_1 \in A(D) \text{ and } w \in V(H) \}.$$

Clearly, D' is strong and the arcs lost in constructing D' are  $x_1x_2$  and those from  $\{x_3, \ldots, x_k\}$  to  $x_1$ . So D' has at least

$$\binom{n-1}{2} + 3 - (k-1) \ge \binom{n-1}{2} + 3 - (n-2) = \binom{n-2}{2} + 3$$

arcs. By the induction hypothesis, we may choose a vertex  $u \in V(D') \setminus \{v^*\}$  such that D'-u is strong. Since  $D' - x_2$  is not strong, we have  $u \neq x_2$ . Moreover, there exists an arc  $wx_2$  with  $w \in V(H)$  and  $w \neq u$ . It is now clear that D - u is also strong.

Next, we suppose that k = 2. Then  $H = D - \{x_1, x_2\}$ . The arcs lost in removing the vertices  $x_1, x_2$  from D are those from V(H) to  $x_1$  and those from  $x_2$  to V(H) apart from  $x_1x_2$ .

Let

 $I_1 = \{wx_1 : w \in V(H) \text{ and } wx_1 \notin A(D)\}$ 

and

$$I_2 = \{ x_2 w : w \in V(H) \text{ and } x_2 w \notin A(D) \}.$$

Suppose that  $|I_1| + |I_2| \ge 2$ . If n = 5, then H is a strong oriented graph with 3 vertices and so H is a directed cycle of length 3. It follows that

$$|A(D)| = |A(H')| + |A(H)| + 2(n-2) - (|I_1| + |I_2|)$$
  

$$\leq 1 + 3 + 2 \times 3 - 2$$
  

$$= 8 < 9$$
  

$$= \binom{n-1}{2} + 3,$$

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a contradiction. This implies that  $n \ge 6$ . So we have that

$$|V(H)| = n - 2 \ge 4$$

and

$$\begin{aligned} |A(H)| &= |A(D)| - |A(H')| - (2(n-2) - (|I_1| + |I_2|)) \\ &\geqslant \binom{n-1}{2} + 3 - 1 - (2(n-2) - 2) \\ &= \binom{n-3}{2} + 3. \end{aligned}$$

Apply the induction hypothesis to H, we can find a vertex u of H such that  $u \neq v^*$  and H-u is strong. If D-u is strong, we are done. If D-u is not strong, then u is either the sole in-neighbor of  $x_1$  in H or the sole out-neighbor of  $x_2$  in H, but not both. In fact, if u is the sole in-neighbor of  $x_1$  in H and is the sole out-neighbor of  $x_2$  in H, then

$$|A(H)| = |A(D)| - 3 \ge \binom{n-1}{2} > \binom{n-2}{2}$$

contradicting the fact that D is an oriented graph. So, by duality, we may assume that u is the sole out-neighbor of  $x_2$ , but not the sole in-neighbor of  $x_1$ . Let  $A_1 = \{x_2w : uw \in A(D) \text{ and } w \in V(H)\}$  and let  $D' = (D - u) + A_1$ . Clearly, D' is strong and the possible arcs lost in constructing D' are  $ux_1, x_2u$  and those from V(H) to u. Since H is a strong digraph, there exists at least one out-neighbor v of u in V(H). Combining this with the fact that H is an oriented graph, we have  $vu \notin A(D)$ . So D' has at least

$$\binom{n-1}{2} + 3 - (2 + (n-4)) = \binom{n-2}{2} + 3$$

arcs. By the induction hypothesis, we may choose a vertex  $u' \in V(D') \setminus \{v^*\}$  such that D' - u' is strong. Clearly,  $u' \neq x_2$ . It follows that D - u' is strong.

Suppose that  $|I_1| + |I_2| \leq 1$ . Then, by duality, we may assume that  $x_1$  is dominated by every vertex of H, and  $x_2$  dominates every vertex of H but at most one vertex. If  $x_2$ dominates  $v^*$ , let T be a spanning out-arborescence of H rooted at  $v^*$  and let u be a leaf of T. Then  $u \neq v^*$  and the subdigraph D - u is strong. If  $x_2$  does not dominate  $v^*$ , then  $x_2$  dominates every vertex of  $V(H) \setminus \{v^*\}$ . Choose an arc  $wv^* \in A(H)$  and let u be an arbitrary vertex in  $V(H) \setminus \{w, v^*\}$ . Then D - u is strong. The proof is complete.

Next we introduce a class of oriented graphs  $D_n$  that will show that the bound on the number of arcs in Theorem 2.1 is sharp.

**Example 2.1** For an integer  $n \ge 4$ , let  $D_n$  be the oriented graph obtained by starting with a directed cycle  $x_1x_2...x_nx_1$  and adding a set of arcs  $\{x_jx_i : 1 \le i \le j-2 \le n-3\}$  to the cycle. The oriented graph  $D_6$  is shown in Figure 1. The oriented graph  $D_n$  has  $\binom{n-1}{2} + 2$  arcs, is strong, and  $x_n$  is the unique vertex whose removal does not destroy strongly connectivity.



Figure 1. The oriented graph  $D_6$ 

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